

# Tracking and Equalizing Channels in Communication Systems and in Data Compression

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June 28, 2007

## Outline

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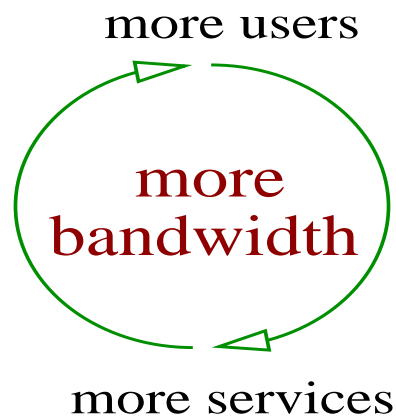
- Training design for estimating wideband wireless channels with antenna arrays at transmitter and receiver
  - ▶ Estimation theoretic perspective
  - ▶ Information theoretic perspective
- Distributed data compression for sensor networks

**Unifying theme:** Channel equalization

## Motivation

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- The increasing demand for ubiquitous wireless access generates a spiraling demand for bandwidth.
- Developing bandwidth efficient physical layer designs is critical.
- The designs need to be optimized for packet transmission.
- Advances in coding and modulation enable us to approach transmission limits, for a given communication channel.
- The challenging issue in wireless communication is channel estimation.

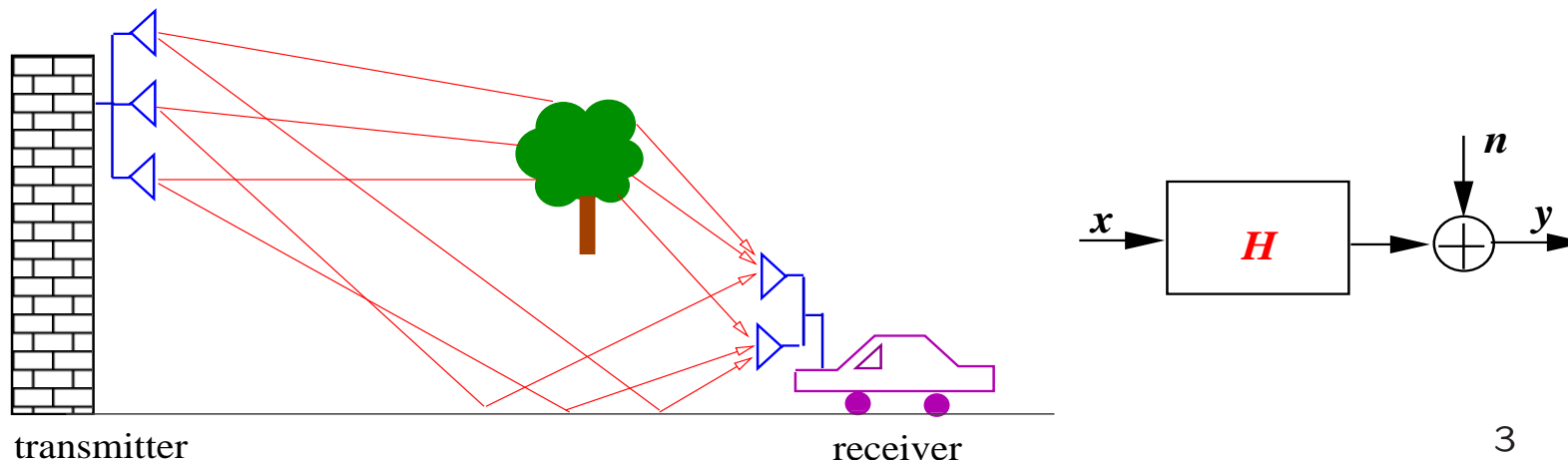


## Channel Estimation in Wireless Communication Systems

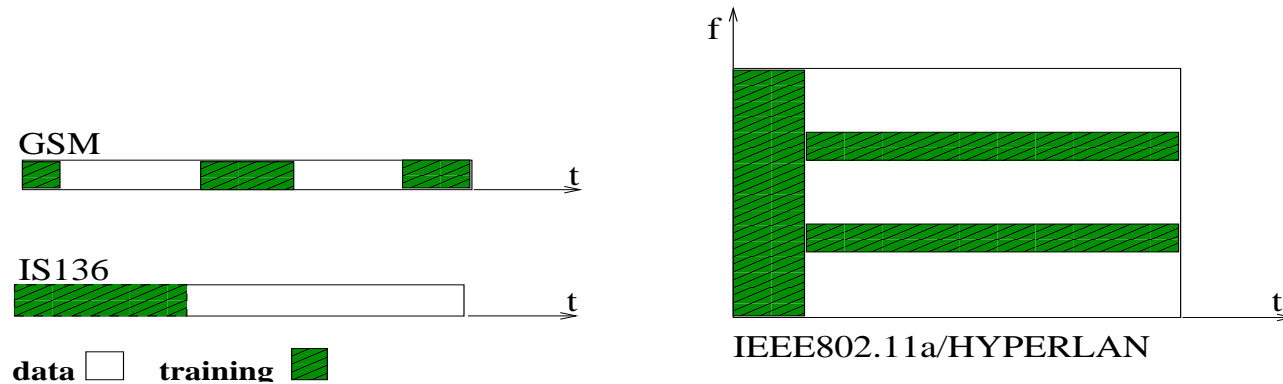
- **Unique characteristics:** fading caused by multipath effect.
- The wireless channel is modelled as a **linear filter**.
- Discrete time representation of the channel input/output relationship:

$$y_i = \mathbf{H}x_i + n_i$$

- To enable channel estimation transmitter multiplexes data symbols  $s$  with known **training** symbols  $t \Rightarrow x$  contains  $s$  and  $t$ .
- Training symbols and the specific multiplexing scheme are known *a priori* at the receiver.



## Training Design in Wireless Communication Systems

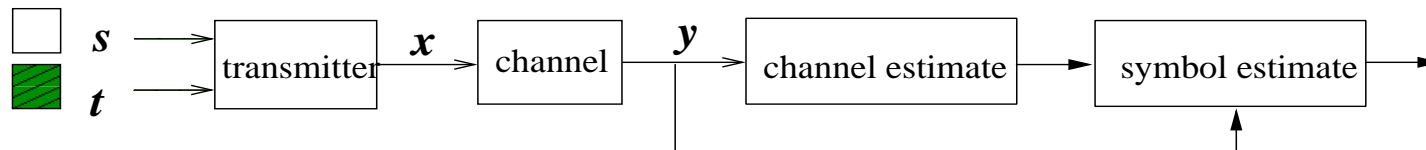


- Current training designs are heuristic.
- Orthogonality between training and data decouples channel and symbol estimation.
- **Tradeoff:** Training improves the system performance; however, it spends the resources that could have been used for transmitting data.
- Investigating optimal training designs is important.
- **Related Work:** literature considered optimal training designs, assuming training and data symbols are separated either in time or frequency.

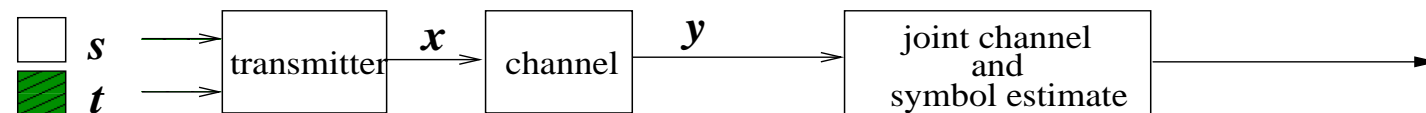
## Sequential versus Joint Channel and Symbol Estimation

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### Sequential Estimate



### Joint Estimate

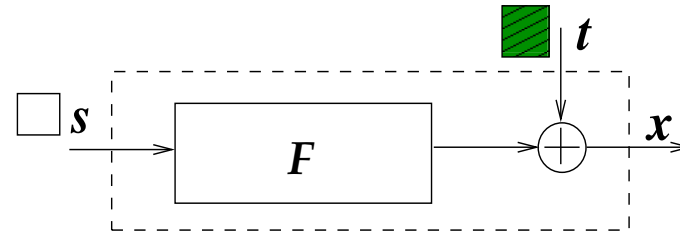


### Questions

- Is separating training and data symbols for sequential estimation optimal?
- How is optimal training design for sequential estimation different from that of joint estimation?

## Affine Precoding: A Unified Transmission Model

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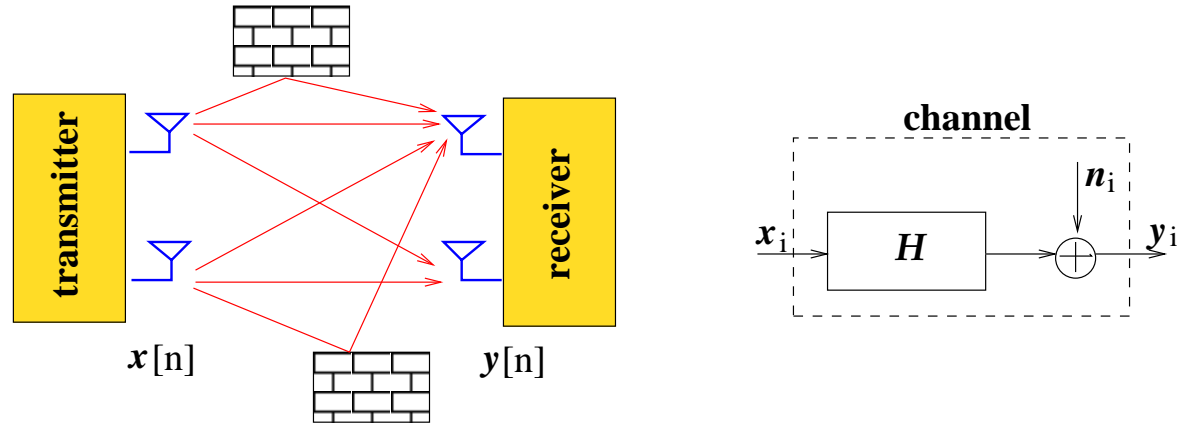
Affine Precoding

- $s \sim \mathcal{CN}(\mathbf{0}, \sigma_{ss}^2 \mathbf{I})$  is a block of user's data and  $t$  is training vector.
- **Affine precoding**:  $x$  is linear precoded data  $Fs$  superimposed on  $t$ .

$$x = Fs + t$$

- Power constraint on transmit vector:  $\mathcal{P} = E\{\|x\|^2\} = \mathcal{P}_s + \mathcal{P}_t$ .
- Affine precoding includes several multiplexing techniques:
  - ▶ GSM and IS136: columns of  $F$  are unit vectors.
  - ▶ IEEE 802.11a/HYPERLAN:  $F$  is FFT matrix.

## Channel Model: Block Fading frequency Selective MIMO



$$y_i = Hx_i + n_i$$

### Assumptions

- ▶  $K$  transmit and  $R$  receive antennas,
- ▶ Length of a transmission block is  $P$ , length of channel memory is  $L$
- ▶ Transmit vector  $\mathbf{x}_i = \text{vec}(\mathbf{x}[iP], \mathbf{x}[iP + 1], \dots, \mathbf{x}[iP + P - 1])$ ,
- ▶ Receive vector  $\mathbf{y}_i = \text{vec}(\mathbf{y}[iP + L], \mathbf{y}[iP + L + 1], \dots, \mathbf{y}[iP + P - 1])$ ,
- ▶ Channel vector  $\mathbf{h} = \text{vec}([\mathbf{H}[0] \dots \mathbf{H}[L]]^T)$ ,
- ▶ Rayleigh fading  $\mathbf{h} \sim \mathcal{CN}(\mathbf{0}, \mathbf{R}_h)$ ,  $n_i \sim \mathcal{CN}(\mathbf{0}, \sigma_{nn}^2 \mathbf{I})$ ,  $n_i$ , and  $\mathbf{h}$  are independent.



## Two Measures of System Performance

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- **Minimizing MSE of channel estimation:**
  - We consider an estimation theoretic bound (Cramer-Rao Bound).
  - CRB establishes a lower bound on the MSE of any estimate of  $\mathbf{h}$ :

$$E\{(\mathbf{h} - \hat{\mathbf{h}})(\mathbf{h} - \hat{\mathbf{h}})^H\} \geq \mathcal{C}_h$$

- Optimal training is the one that minimizes  $\text{trace}(\mathcal{C}_h)$ , where  $\mathcal{C}_h$  is CRB for estimating channel.
- **Maximizing information rate:**
  - We consider a lower bound on mutual information  $I(\mathbf{s}; \hat{\mathbf{s}})$ .
  - Most receivers make the decision about the symbols based on **soft estimate of the modulated symbols** and  $I(\mathbf{s}; \hat{\mathbf{s}}) \leq I(\mathbf{s}; \mathbf{y})$ .
  - $I(\mathbf{s}; \hat{\mathbf{s}})$  measures the reduction in the uncertainty about  $\mathbf{s}$  due to the knowledge of  $\hat{\mathbf{s}}$ .
  - Optimal training is the one that maximizes this lower bound.

## Minimizing Channel Estimation Error

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**Notation**  $y = Hx + n = H(Fs + t) + n \Rightarrow Ht = \Phi(t)h$  and  $Hf_i = \Phi(f_i)h$ .

- **Optimality condition for sequential estimation**

Minimizing  $tr(C_h)$  requires:

$$\Phi(t)^H \Phi(f_i) = 0 \quad i = 1, 2, \dots, N$$

► **Remark 1:** To satisfy orthogonality constraint it is sufficient to separate data and training in frequency domain.

► **Remark 2:** Under orthogonality constraint  $C_h = (\sigma_{nn}^{-2} \Phi(t)^H \Phi(t) + E_h\{\Xi\} + R_h^{-1})^{-1}$ .

- **Optimality condition for joint estimation**

Minimizing  $trace(C_h)$  requires:

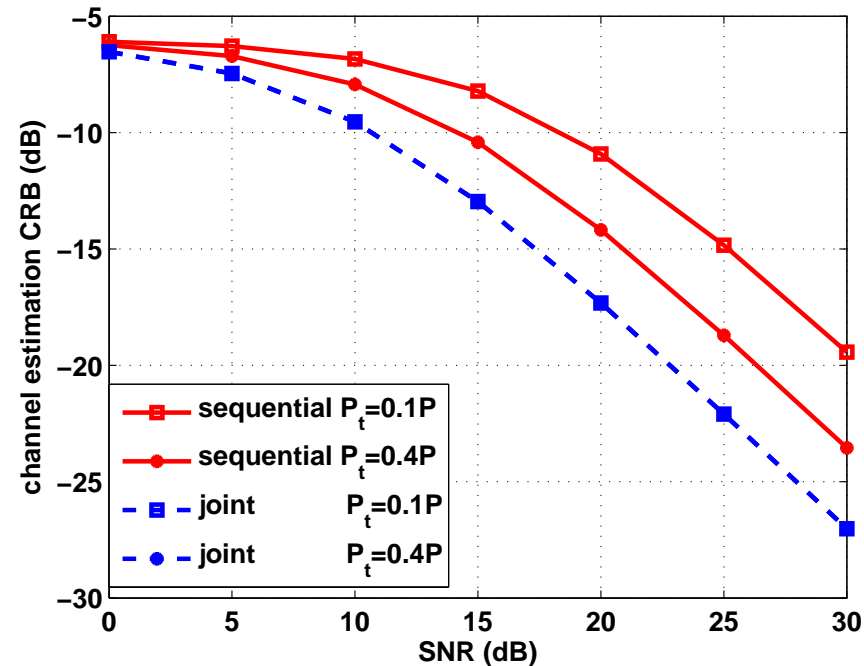
$$\sigma_{ss}^2 \sum_{i=1}^N \Phi(f_i)^H \Phi(f_i) + \Phi(t)^H \Phi(t) = \Lambda$$

where  $\Lambda$  is diagonal.

► **Remark 3:** For diagonalization no orthogonality constraint is required to be satisfied, and  $C_h = (\sigma_{nn}^{-2} \Lambda + R_h^{-1})^{-1}$ .

## Power Allocation Minimizing Channel Estimation Error

- $K = R = 2, L = 3, \sigma_{ss}^2 = 1$



- For sequential estimate minimizing  $trace(\mathcal{C}_h)$  requires  $\mathcal{P} = \mathcal{P}_t$ . For joint estimate minimum of  $trace(\mathcal{C}_h)$  depends on  $\mathcal{P}$ , not  $\mathcal{P}_t$ .
- Training design cannot be only based on minimizing channel estimation error  $\Rightarrow$  we consider mutual information  $I(\mathbf{s}; \hat{\mathbf{s}})$ .

## A Lower Bound on Mutual Information $I(\mathbf{s}; \hat{\mathbf{s}})$

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- We link the lower bound on  $I(\mathbf{s}; \hat{\mathbf{s}})$  with symbol estimation performance.

Let  $\hat{\mathbf{s}}$  be any symbol estimate, we have  $E\{(\mathbf{s} - \hat{\mathbf{s}})(\mathbf{s} - \hat{\mathbf{s}})^H\} \geq \mathcal{C}_s$ , where  $\mathcal{C}_s$  is CRB for estimating the symbols.  $I(\mathbf{s}; \hat{\mathbf{s}})$  can be lower bounded as:

$$I(\mathbf{s}; \hat{\mathbf{s}}) \geq \log(|\sigma_s^2 \mathcal{C}_s^{-1}|)$$

- ▶ **Remark 1:** Maximizing lower bound is equivalent to minimizing  $|\mathcal{C}_s|$ .
- ▶ **Remark 2:** A better symbol estimate results in a higher transmission rate.

## Mutual Information Lower Bound For Gaussian Symbols

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- **Lower bound for sequential estimation**

Receiver models  $\mathbf{h} = \hat{\mathbf{h}} + \tilde{\mathbf{h}}$ . Assuming  $\hat{\mathbf{h}}$  is the LMMSE channel estimate:

$$\begin{aligned} I(\mathbf{s}; \hat{\mathbf{s}}) &\geq E_{\hat{\mathbf{h}}} \{ I(\mathbf{s}; \hat{\mathbf{s}} | \hat{\mathbf{h}}) \} \\ &= E_{\hat{\mathbf{h}}} \{ \log(|\sigma_{ss}^2 E_{\mathbf{s}} \{ \mathbf{F}^H \hat{\mathbf{H}}^H \mathbf{R}_{v|s}^{-1} \hat{\mathbf{H}} \mathbf{F} + \boldsymbol{\varepsilon}^H (\mathbf{R}_{v|s}^{-1} \otimes \mathbf{D}^H \mathbf{R}_{v|s}^{-1} \mathbf{D}) \boldsymbol{\varepsilon} \} + I|) \} \end{aligned}$$

► **Remark 1:** bound depends of  $\hat{\mathbf{H}}$  and estimation error covariance.

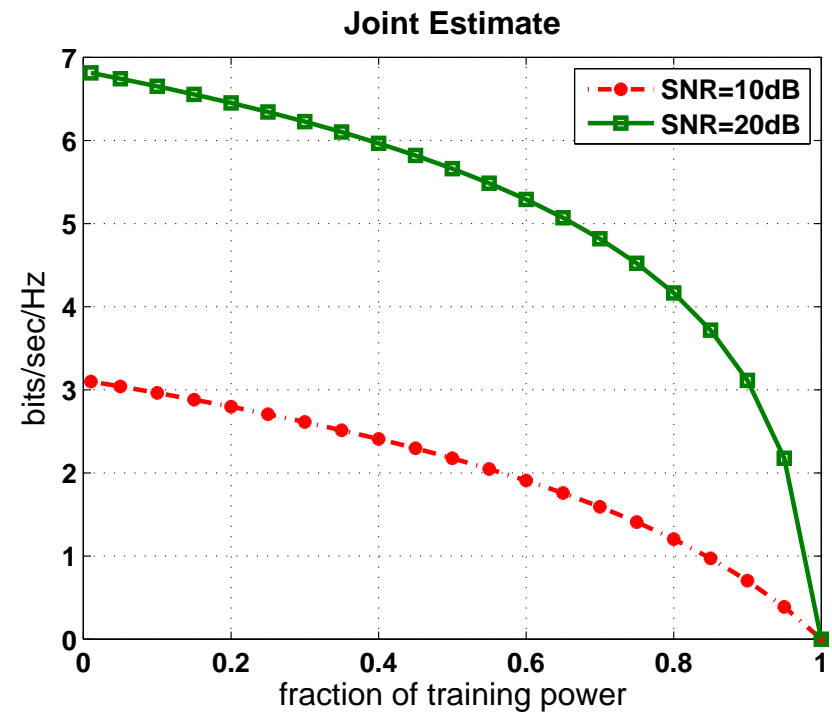
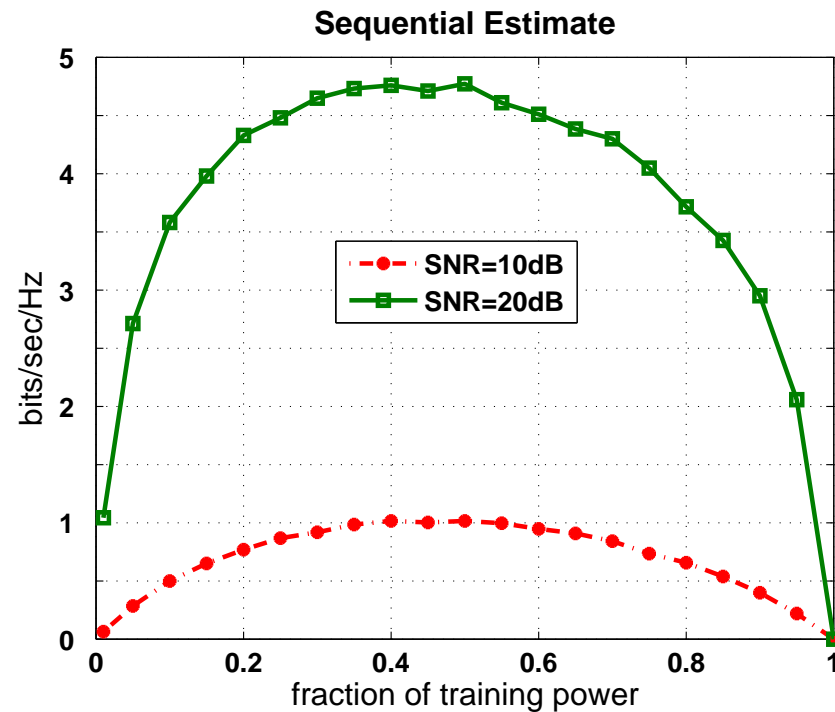
- **Lower bound for joint estimation**

$$\begin{aligned} I(\mathbf{s}; \hat{\mathbf{s}}) &\geq \log(|\sigma_{ss}^2 \mathbf{C}_s^{-1}|) \\ \mathbf{C}_s^{-1} &= \sigma_{nn}^{-2} \begin{bmatrix} \text{tr}(\mathbf{R}_h \mathbf{F}_1^H \mathbf{F}_1) & \cdots & \text{tr}(\mathbf{R}_h \mathbf{F}_1^H \mathbf{F}_N) \\ \vdots & \ddots & \vdots \\ \text{tr}(\mathbf{R}_h \mathbf{F}_N^H \mathbf{F}_1) & \cdots & \text{tr}(\mathbf{R}_h \mathbf{F}_N^H \mathbf{F}_N) \end{bmatrix} + \sigma_{ss}^{-2} \mathbf{I} \end{aligned}$$

► **Remark 2:** bound is independent of  $\hat{\mathbf{H}}$  and estimation error covariance.

► **Remark 3:** maximizing bound leads to the same optimality conditions derived for minimizing channel estimation error.

## Power Allocation Maximizing Mutual Information Lower Bound



- Minimizing the bound requires  $\mathcal{P}_t \simeq 0.4\mathcal{P}$  for sequential estimate and  $\mathcal{P}_t = 0$  for joint estimate.

## Summary

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- We investigated optimal training for transmission over wideband wireless MIMO channels.
- We adopted affine precoding and considered channel estimation error and information rate as figures of merit.
- We showed that optimal training design depends on the receiver structure.
- For joint estimate training is not needed.
- Implementing joint estimate is computationally expensive.
- For sequential estimation training is needed for minimizing channel estimation error and maximizing mutual information lower.
- For sequential estimation separating data and training in frequency is optimal.

## Future Work

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- Optimal affine precoding design for practical communication systems:
  - ▶ specific channel and symbol estimation algorithms,
  - ▶ practical measures for system performance (e.g., BER),
  - ▶ specific channel coding and modulation schemes.
- Tradeoff between training and error correction channel coding.



## Outline

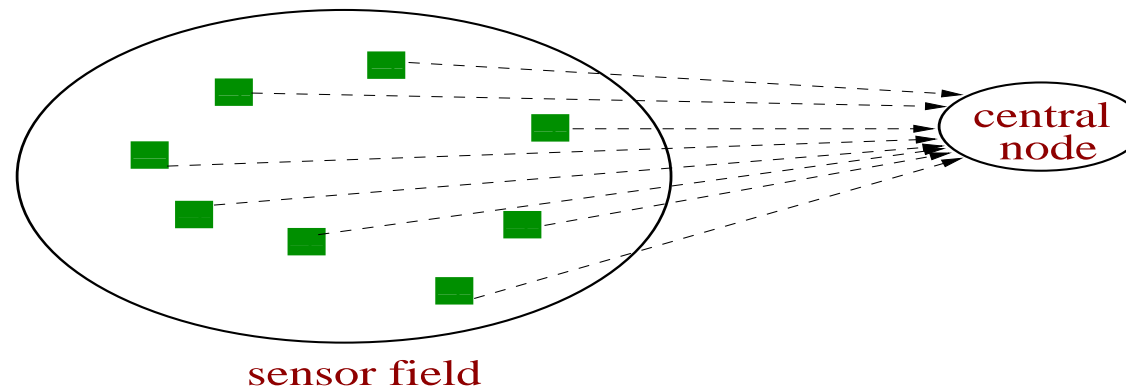
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- Training design for estimating wideband wireless channels with antenna arrays at transmitter and receiver
  - ▶ Estimation theoretic perspective
  - ▶ Information theoretic perspective
- Distributed data compression for sensor networks

## Motivation

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- Sensor networks consist of many low power and cheap sensors with limited computational capabilities.
- Neighboring sensors measure **highly correlated** data.
- Sensors send their measured data to a central node for joint decoding.
- The sensors could remove redundant data, if joint encoding is allowed.
- **Question:** is there a way of removing the redundancy in a completely **distributed** manner (without requiring the sensors to inter-communicate)?

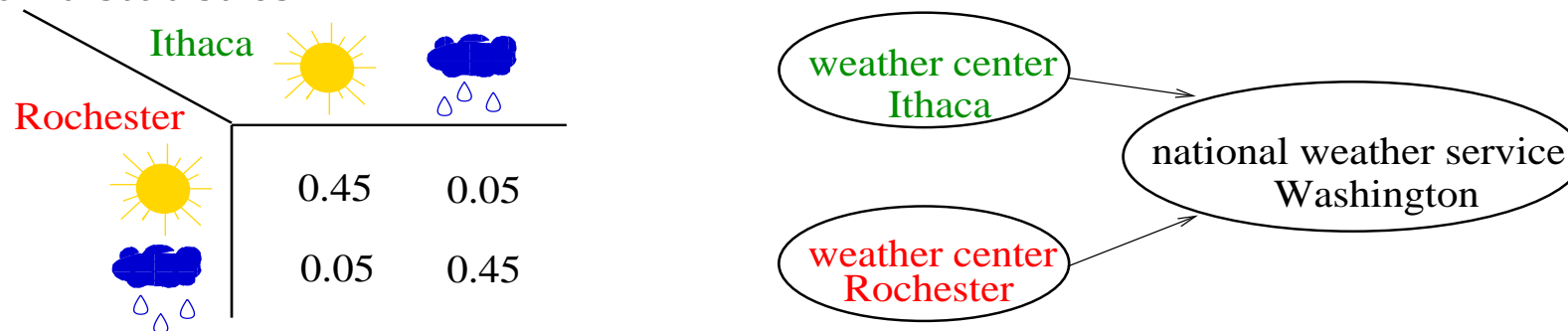


## Coding with Side Information: Example

- Coding with side information enables distributed data compression.

### Example

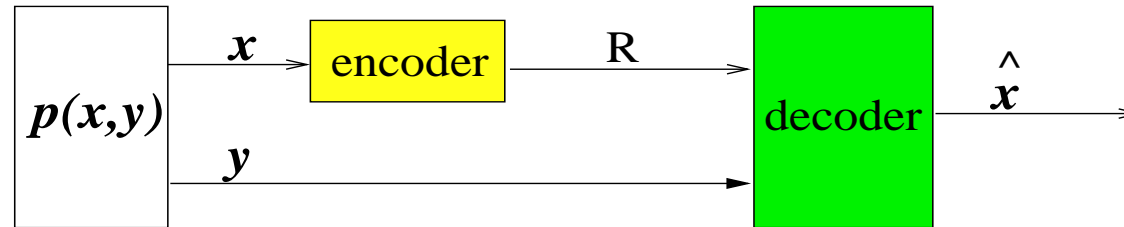
- ▶ Rochester and Ithaca report the weather of 100 days in Summer to Washington.
- ▶ Washington knows the weather at Rochester, and the joint statistics.
- ▶ Ithaca is not aware of the weather in Rochester; however, it knows the joint statistics.



### Question

- ▶ How many bits are needed to convey the Ithaca weather to Washington?

## Coding with Side Information: Model



**Set up:**  $x, y$  are related through  $p(x, y)$ , known at encoder and decoder.

- **(Slepian and Wolf coding)** for **discrete**  $x, y$  one can compress  $x$  into  $R < R_x$  bits to reconstruct  $x$  perfectly.

▶ **Answer to the question:** for independent coding  $R_x = 100$  bits, and for coding with side information  $R = 50$  bits are needed.

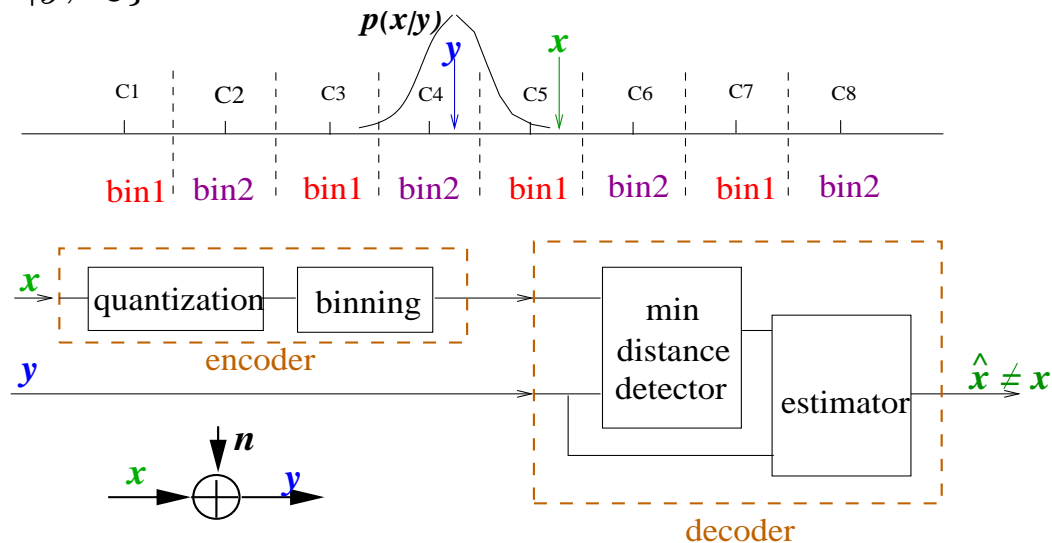
- **(Wyner and Ziv coding)** for **continuous**  $x, y$  one can compress  $x$  into  $R < R_x$  bits to reconstruct  $x$  with MSE distortion  $D$  (e.g.,  $\mathbb{E}\{(x - \hat{x})^2\} = D$ )

**Conclusion:** exploiting correlation at encoder/decoder allows a reduction of transmission rate.

**Challenge:** proofs are not constructive and are based on the idea of **binning**.

## Most Relevant Work: DISCUS [Pradhan and Ramchandran 2003]

- DISCUS is a **constructive** approach to the problem of coding with side information.
- **Assumption:** the sources  $x, y$  are related by  $y = x + n$ .
  - ▶ **Encoder** quantizes  $x$  to  $c_5$ , and finds the bin index  $p = 1$ .
  - ▶ Given  $p$  and  $y$  **decoder** decides  $c_5$  based on minimum distance rule, and forms  $\hat{x} = \mathbb{E}\{x|y, c_5\}$ .

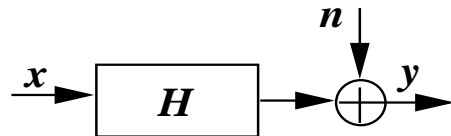


- **Question:** What if the relationship between  $y$  and  $x$  **cannot** be modelled as  $y = x + n$ ?

## A General Linear Correlation Model

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- We consider two vector sources  $\mathbf{x}, \mathbf{y} \in \mathbb{R}^N$  that are related by

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{n}$$


### Assumptions

- ▶  $\mathbf{H}$  is constant,  $\mathbb{E}\{\mathbf{x}\mathbf{x}^T\} = \mathbf{R}_{xx}$ ,  $\mathbb{E}\{\mathbf{n}\mathbf{n}^T\} = \mathbf{R}_{nn}$ ,  $\mathbf{x}$  and  $\mathbf{n}$  are independent,
- ▶ Decoder knows  $\mathbf{y}$  and the parameters  $\mathbf{H}$ ,  $\mathbf{R}_{xx}$ ,  $\mathbf{R}_{nn}$ ,
- ▶ Encoder knows neither  $\mathbf{y}$  nor  $\mathbf{H}$ ,  $\mathbf{R}_{nn}$ .

**Goal:** develop low complexity compression techniques to encode  $\mathbf{x}$ .

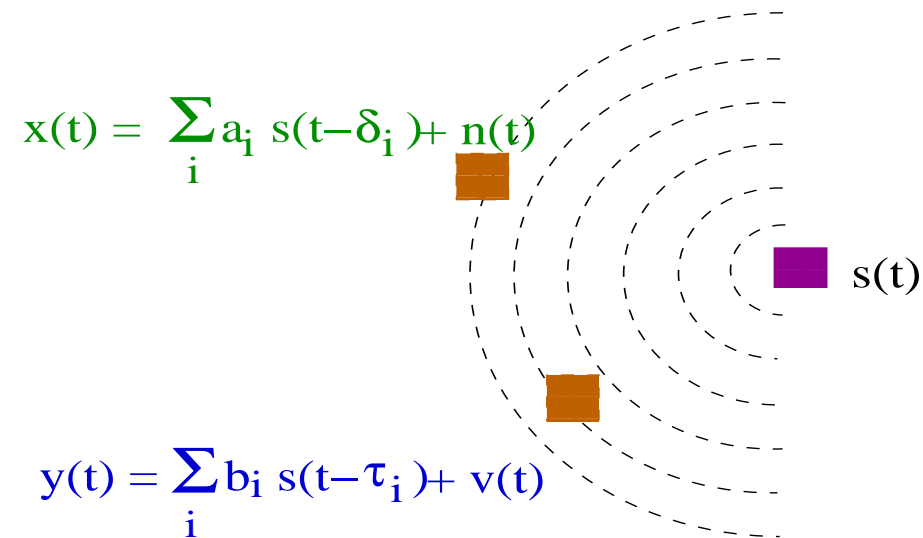
## Why a Linear Correlation Model?

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$$y = Hx + n$$

- It is more general than the model  $y = x + n$ .

Potential Application: audio field sensors



- For jointly Gaussian sources the model is exact.

## Relevant Questions

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- Can we use the code construction designed for scalar source coding (e.g., DISCUS) to solve vector source coding problem?
- What is the structure of encoder/decoder?
- What is the best rate allocation to encode components of  $x$ , subject to a sum rate constraint?

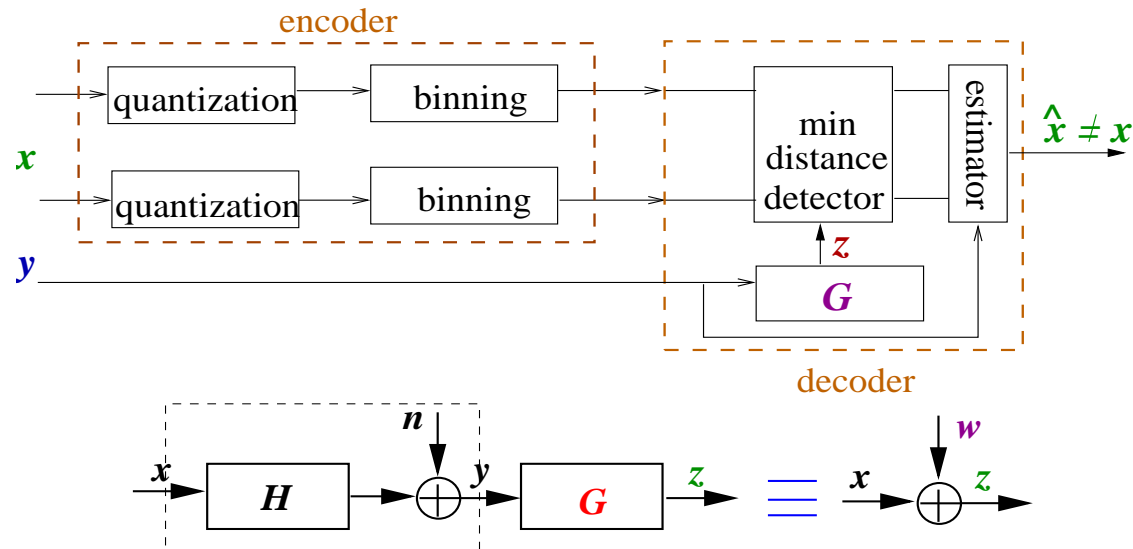


## Encoder and Decoder Architecture: Equalization at the Decoder

- We use tools from communications (linear equalizer  $G$ ) to convert vector source coding problem into several scalar source coding problems.

$$y = Hx + n \Rightarrow z = Gy = x + w$$

Zero Forcing  $G : GH = I$  LMMSE equalizer  $G = \arg \min \mathbb{E}\{\|Gy - x\|^2\}$



- Given compression rate  $r_i$  bits/sample encoder quantizes  $x_i$  and provides decoder with bin indices  $p_i$ . Given  $p_i$  and  $z_i$  decoder decides on quantization intervals, and forms  $\hat{x}$ .

## Rate Allocation Policy in Distributed Data Compression

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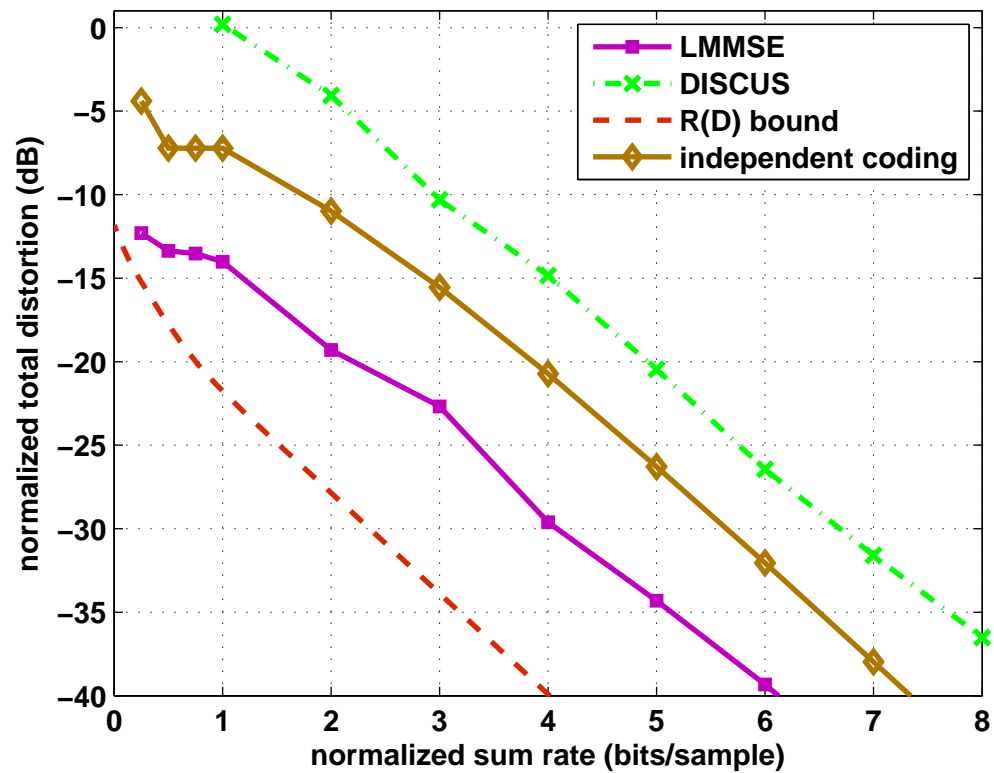
- Performance of the proposed system depends on rate allocation policy.
- Encoder compresses  $x_i$  with  $r_i$  bits/sample.
- **Optimal rate allocation:** given  $R = \sum_{i=1}^N r_i$  what is the best  $\{r_i\}_{i=1}^N$ ?

$$\mathbf{y} = \mathbf{H}\mathbf{x} + \mathbf{n} \Rightarrow \mathbf{G}\mathbf{y} = \mathbf{z} = \mathbf{x} + \mathbf{w}$$

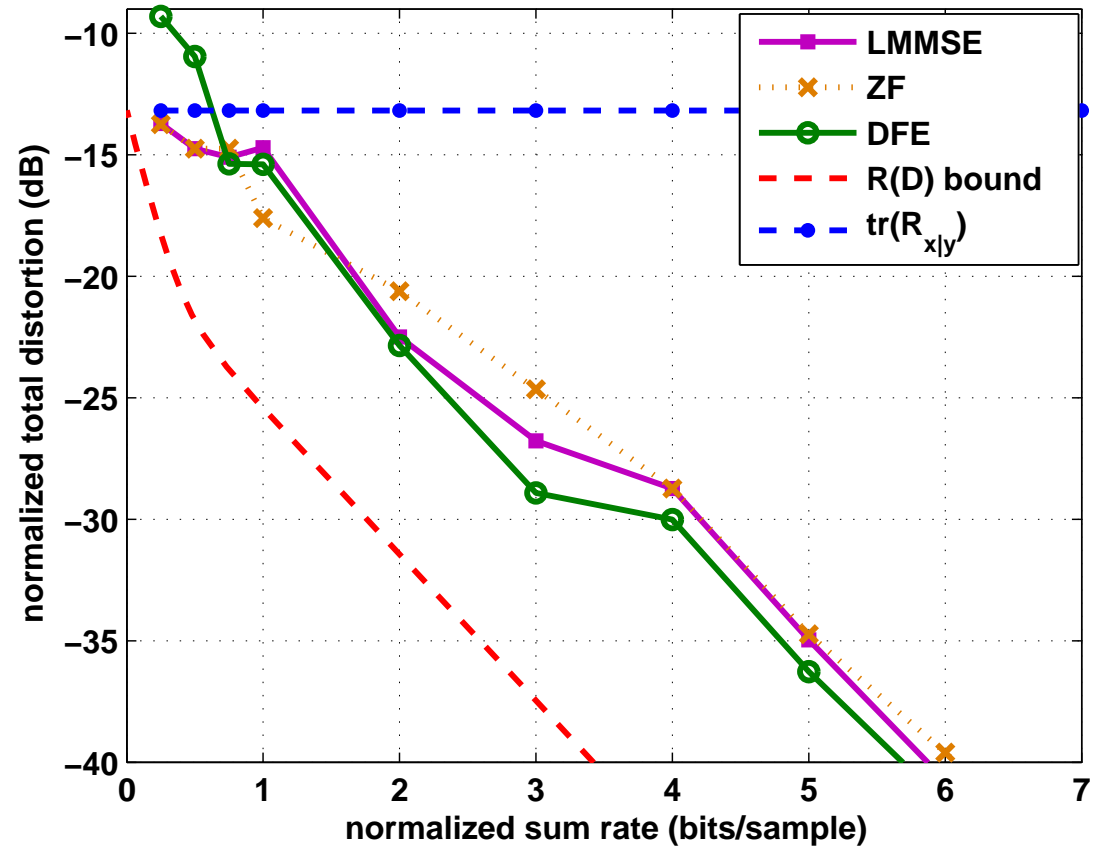
- Statistics of  $\mathbf{w}$  depends on  $\mathbf{G}$ .
- Each  $z_i$  corresponds to a different  $\text{var}(w_i) \Rightarrow r_i$  should be adapted to  $\text{var}(w_i)$ .
- **Proposition:** assign bits such that higher  $\text{var}(w_i) \Rightarrow$  more  $r_i$ .
- ▶ **Remark 1:** the rate allocation policy contradicts the one in communication systems where higher  $\text{var}(w_i) \Rightarrow$  less  $r_i$ .
- ▶ **Remark 2:** higher  $\text{var}(w_i) \Rightarrow z_i$  is less correlated is to  $x_i \Rightarrow$  encoder needs to send more bits to the decoder.

## Empirical Rate Distortion Bound: DISCUS vs. Decoder With LMMSE

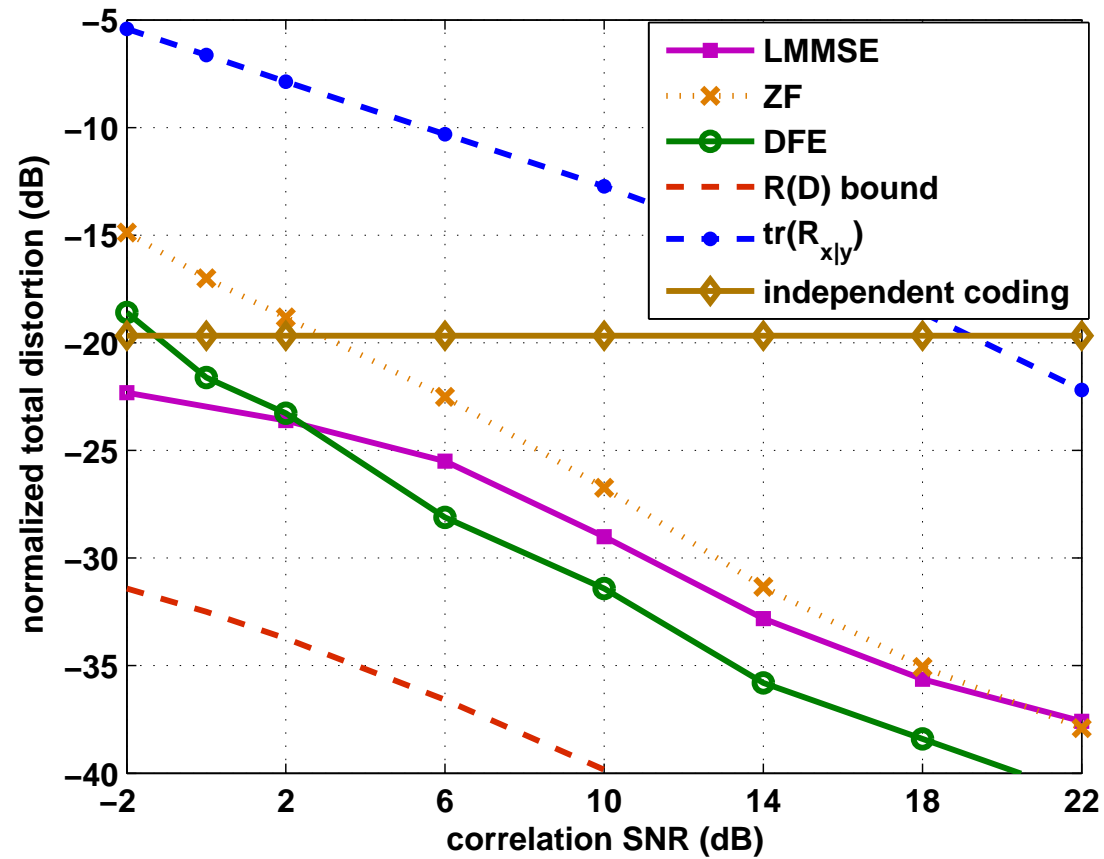
- $N = 4$ ,  $\mathbf{H}$  and  $\mathbf{R}_{xx}$  are Toeplitz,  $\mathbf{x} \sim \mathcal{N}(0, \mathbf{R}_{xx})$  and  $\mathbf{n} \sim \mathcal{N}(0, \sigma_n^2 \mathbf{I})$  with  $10 \log_{10} \sigma_n^2 = -12$  dB.



## Empirical Rate Distortion Bound: Decoder With LMMSE,ZF,DFE



## Distortion vs. Correlation SNR :Decoder With LMMSE,ZF,DFE



## Summary

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- We considered coding of vector source  $x$  with side information  $y$  when the correlation between  $x$  and  $y$  is linear  $y = Hx + n$ .
- Using the equalization techniques we provided low complexity compression algorithms.
- The compression scheme reduces transmission rate, without explicitly assuming the correlation model at the encoder.
- We proposed a rate allocation policy which minimizes error probability of minimum distance detection.

## Future Work

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- How can we obtain parameters of the correlation model at the decoder?
- How can we make the bin index transmission robust to possible noise?
- How can we extend this code construction to more than two vector sources?

