

Issues and Approaches to Management of Sensor Networks

Kang G. Shin

Real-Time Computing Laboratory

Department of EECS

University of Michigan

Outline of Talk

- Generics, issues, and approaches of sensor networks
 - Hardware
 - Communication
 - Software
- Examples of research on sensor networks at:
 - University of Michigan
 - UC Berkeley
 - University of Virginia
 - UCLA
 -

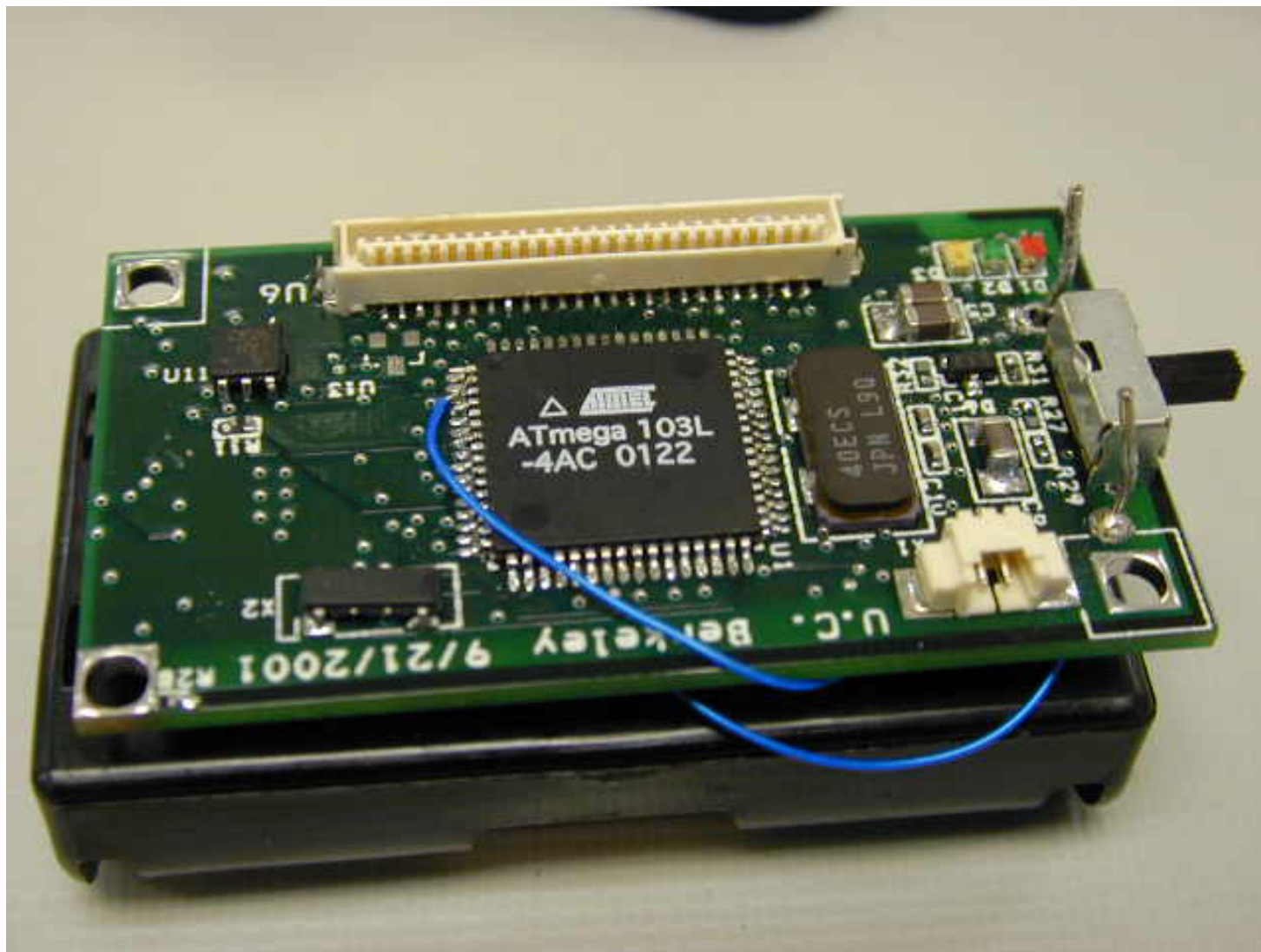
Characteristics of Sensor Networks

- Large # of small, resource-limited sensor nodes, operating *in aggregate*
 - Usually battery-powered, hence *energy-constrained*
 - Wide range of sensing capabilities
 - Temperature, light, sound, magnetic fields, motion, vision
 - Low-power wireless networking
 - Unattended, inaccessible, prolonged deployment
 - Requires *in-network processing*
 - *Time-varying* functions/roles
- => Must be **self-organized**, **self-maintaining** and **programmed *in situ*** to operate at very low duty cycle

Uses of Sensor Networks

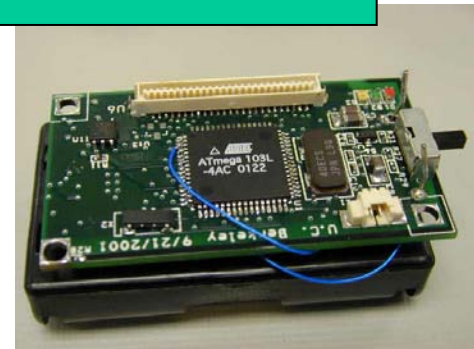
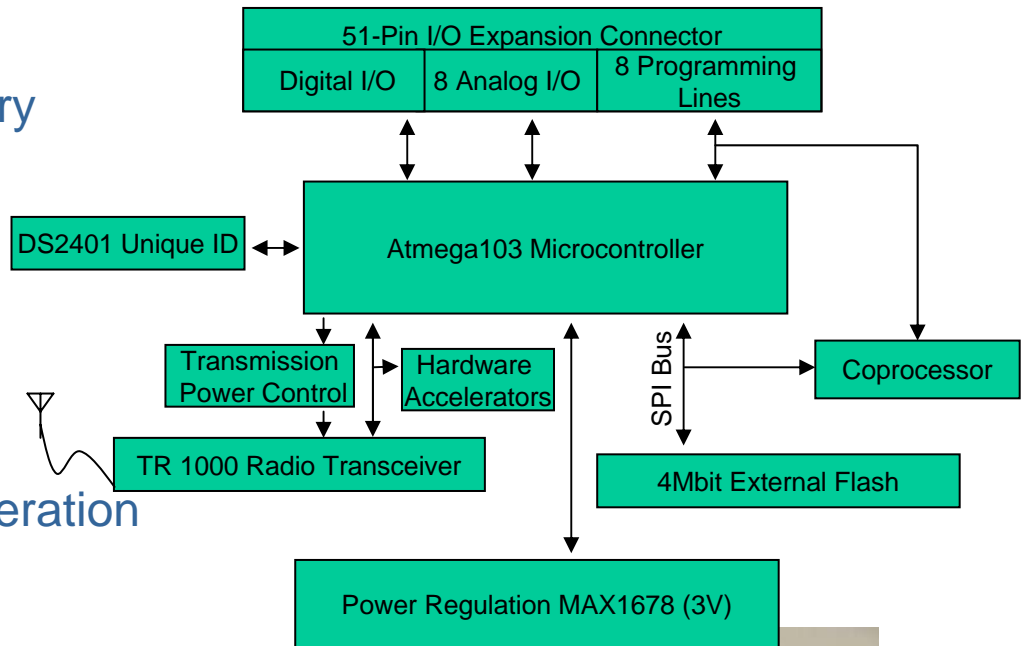
- **Commercial**
 - Manufacturing plant monitoring, integrated robotics, vehicle/object tracking, security/safety monitoring, inventory control and manuals/instructions (RFIDs), etc.
- **Research**
 - Environmental monitoring (habitat and agricultural studies)
- **Military**
 - Tracking, intrusion detection
- **Homeland security**
 - Surveillance of public/critical infrastructures such as buildings, bridges, utility distribution and water supply systems

Typical Sensor Node: X-Bow Mica Mote



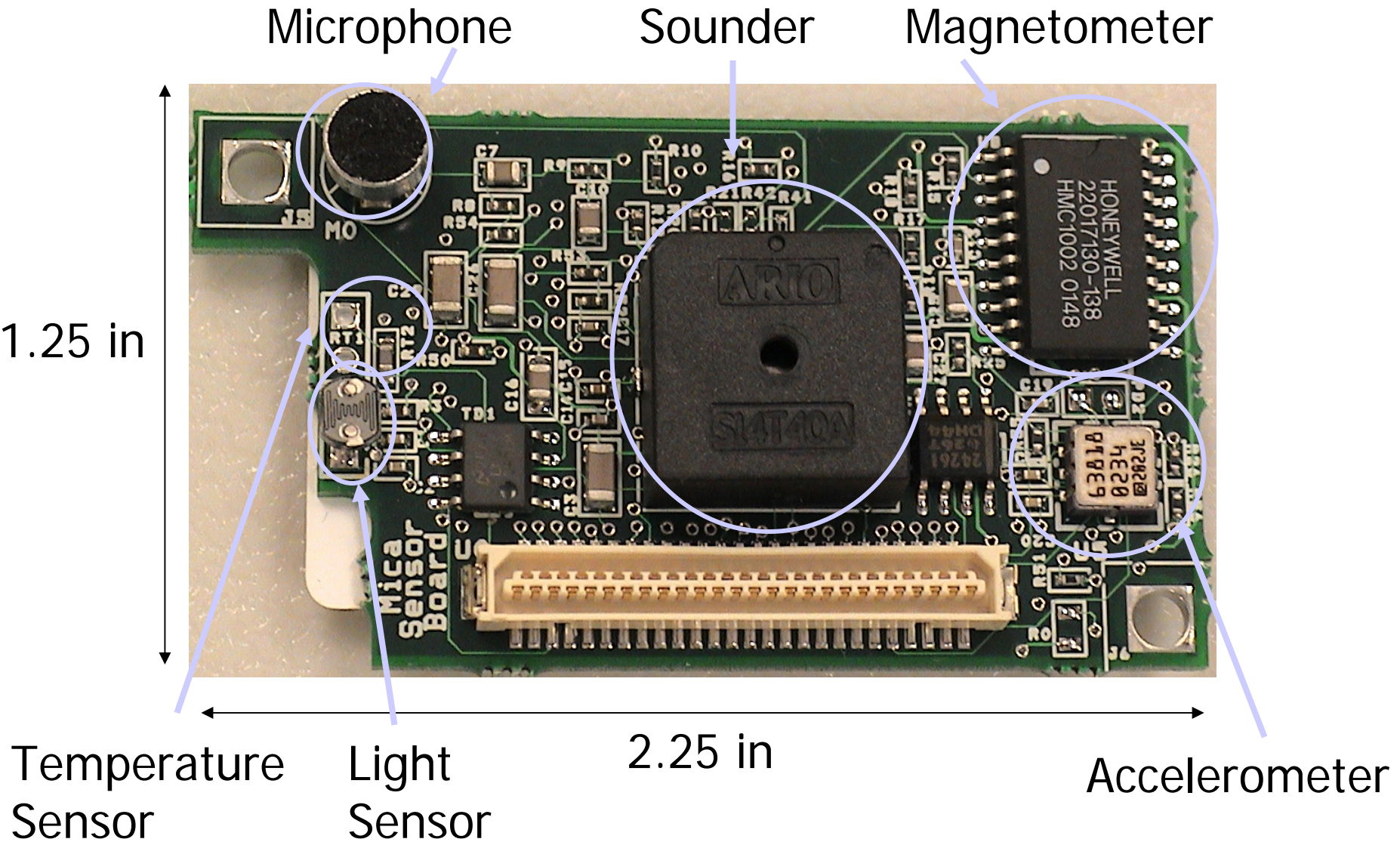
MICA Architecture

- Atmel ATMEGA103
 - 4 MHz 8-bit CPU
 - 128KB Instruction Memory
 - **4KB RAM**
- 512 KB flash (AT45DB041B)
 - SPI interface
 - 1-4 $\mu\text{j/bit}$ r/w
- RFM TR1000 radio
 - 50 kb/s – ASK
 - Focused hardware acceleration
- 6 ADC channels
- **Unique serial IDs**
- Network programming
- 51-pin connector
 - Analog compare + interrupts
- TinyOS tool chain



2xAA form factor Cost-effective power source

Sensor Board Device Placement



Mica Mote with Sensor Board



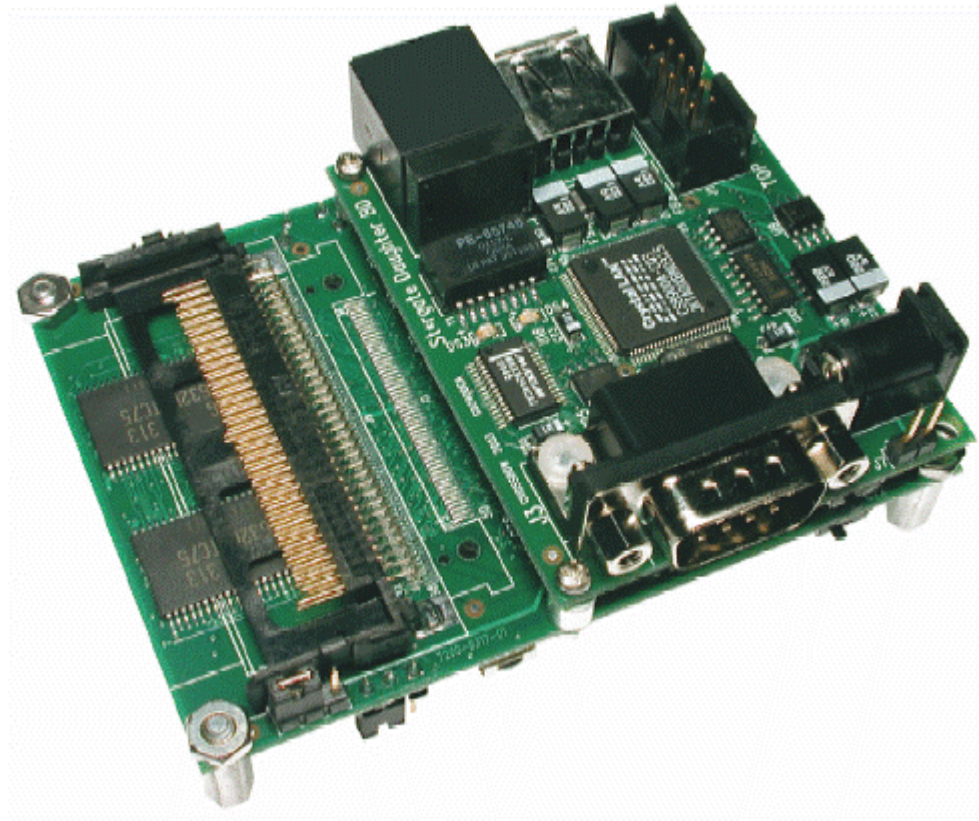
Research Areas/Issues

- **Sensing and architecture**
 - Sensor hardware design (MEMS)
 - Signal/Data processing
 - Rich interfaces and simple primitives allowing cross-layer optimization
 - Low-power processor, ADC, radio, communication, encryption
- **Resource management (operating system)**
 - Limited computational power, memory, code space, electrical power
 - Node computation & communication, and their scheduling
- **Networking and distributed services**
 - Medium Access Control & routing
 - Clock synchronization, localization, and data aggregation
- **Programming**
 - Software component models and middleware
 - Describe global behavior, synthesize local rules that have correct, predictable global behavior
- **Applications**
 - Long-lived, self-maintaining, dense instrumentation of previously unobservable phenomena
 - interaction with a computational environment

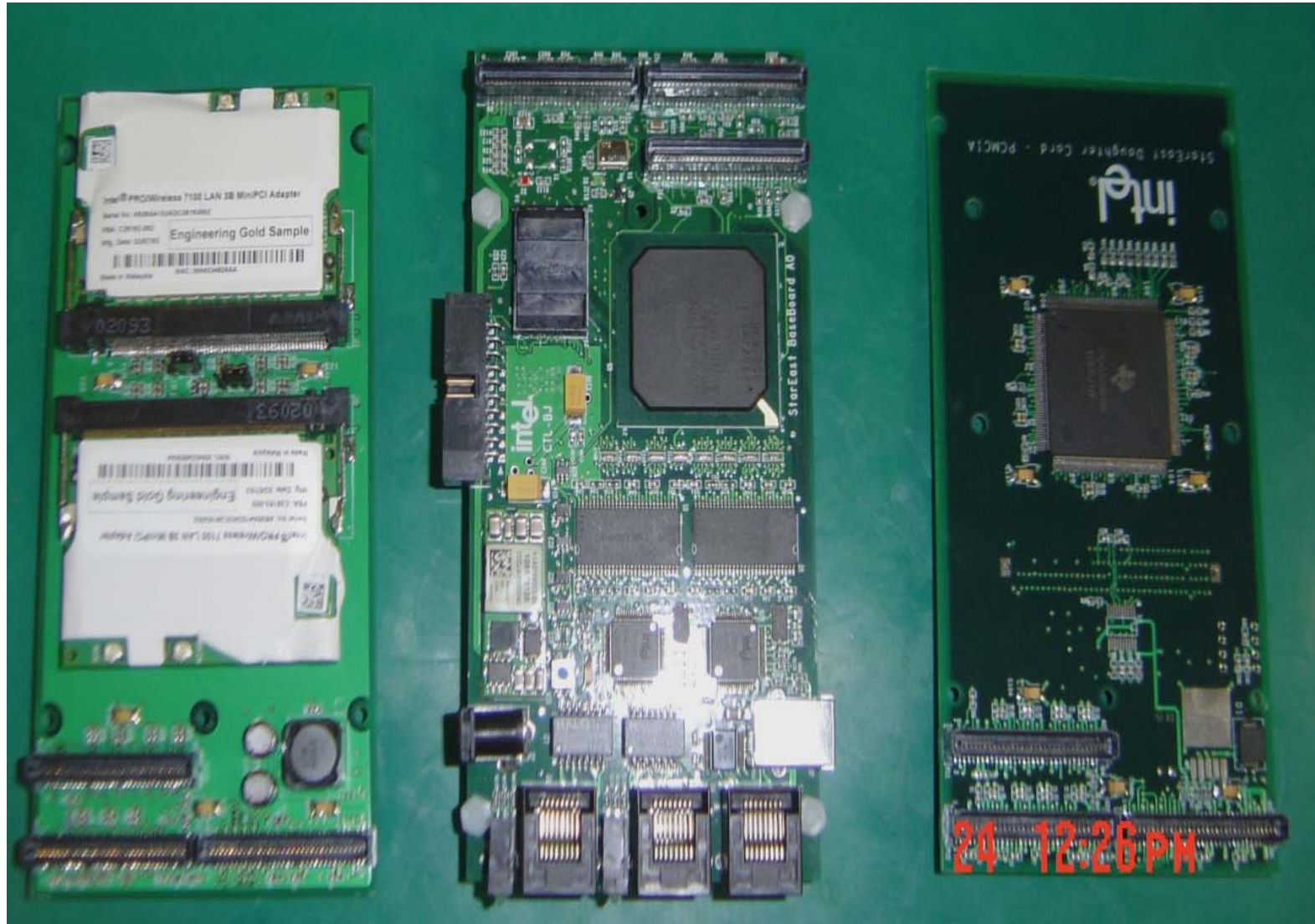
Networking

- **Medium Access Control (MAC)**
 - Main issues with wireless communication
 - Collisions
 - Limited range
 - Hidden node/terminal problem
 - Transmission errors
 - Motes use CSMA (Carrier Sense Multiple Access)
 - Cannot send and receive at the same time
 - Cannot detect collision
 - Work is being done to create inherently collision-free MAC protocols
 - TDMA in a region; may be closely-coupled with applications
 - ... or to reduce the probability of collision
 - **Implicit** acknowledgements
 - S-MAC – coordinates sleep cycles to *save energy* and *avoid collisions*
 - Non-Mote systems (esp. simulations and more powerful sensors) use 802.11 MAC or its variations: Stargate and Stareast

Super Node I: Stargate Board



Super node II: Stareast Boards

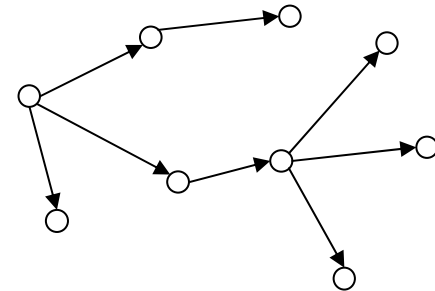


Routing Protocols

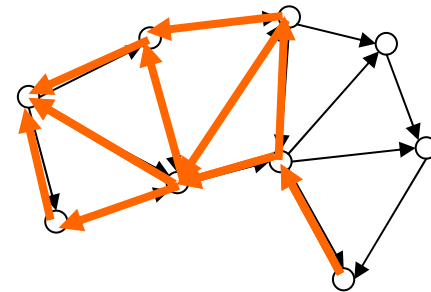
- **Spanning tree** within a cluster/region
- **Geographic routing**
 - Route messages to a specific location
 - Each node knows its location
 - No routing tables maintained
- **Cluster-based routing**
 - Use simple table-based routing protocols to route **to cluster head** (e.g., dynamic source routing, ad-hoc on demand distance vector routing)
 - Use higher-level protocol (e.g., geographic) to route **between** cluster heads
- **Landmark routing**
 - Similar to cluster-based routing, but without the cluster formation overhead
 - Messages are routed to **known landmarks**, from which they are routed to their final destination

Routing Protocols, cont'd

- **Gradient Routing**
 - Requires only **local** information at each node
 - An “interest” is propagated outward by a **sink** node
 - Each node receiving the interest remembers it and passes it along
 - Different topologies arise due to forwarding policies
 - Data from a source traces back links to the sink
 - **Preferred data paths** may be reinforced
 - Lowest energy
 - Shortest path
 - Least latency



Spanning Tree



Directed Acyclic Graph

Sensor Network Programming

- **Embedded systems**
 - Lightweight OS, e.g., tinyOS, EMERALDS
 - OS and application software are compiled and linked together, then downloaded to the node
 - Programmed once and deployed
 - Some work is being done on **network reprogramming**
 - Expensive in terms of energy
 - Takes a node out of service while reprogramming
 - Scalability issues
- **Software structured using component models**
 - Support modularity
 - Only essential components are compiled into the system
 - Easy to upgrade/replace components *during development*

Example University Research Efforts

- University of Michigan
- UC Berkeley
- University of Virginia
- UCLA
-

Efforts at UMICH

- DARPA:
 - SMILE: Service Models for Integration of real-time Embedded systems
 - Security Tradeoffs (with UMass and ASU)
- ONR, NRL, NSF, Cisco:
 - LiSP (Lightweight Security Protocol), PIV, DKMP, SyKeeper
- NSF:
 - Lightweight and Flexible Sensor Network Management
- Project personnel: 1 faculty, 1 full-time research scientist, and 9 grad students
- Project URLs:
<http://kabru.eecs.umich.edu/{smile,security}>

Sensor Network Testbed



Sample Projects at UM

- *Adaptive Query Processing (AQP)*
- *Content-aware metadata creation in a heterogeneous mobile environment*
- *Network routing*
- *Distributed location service*
- *Sensor network security*
- Self-management

AQP Middleware

- Provides an **abstraction** that forms the basis for *service & application development* on a platform
=>Higher-level domain services are implemented as **queries** and **query-triggered functions**
- Is based on a **data-centric view** of networked embedded systems
- Provides **basic** data access and management
- Is based on a **data model** that includes type, time, location, and quality parameters

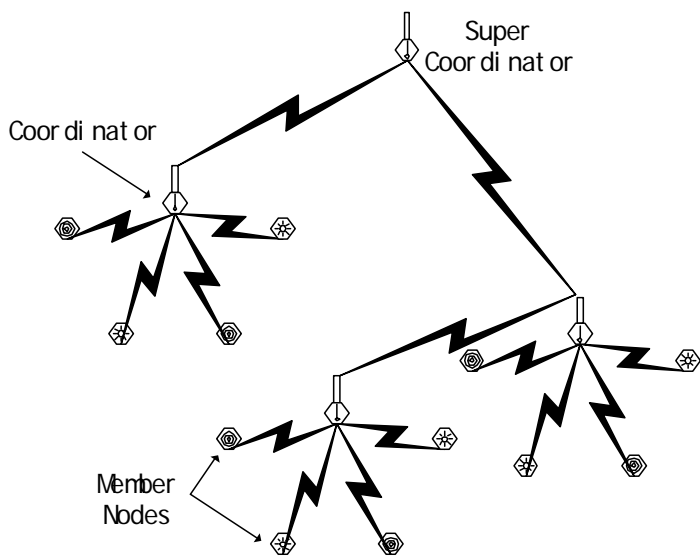
Service Development on Motes

- Sensor database (SensorDB)
- Energy-aware Query Processing
 - Declarative Query Interface to
 - provide transparent **adaptation** and **optimization**
 - Energy savings in
 - communication and query processing
- Techniques proposed to increase lifetime
 - **Utility/cost** in *query allocation* by each coordinator
 - **Energy-efficient** (i.e., computationally-efficient) *query indexing* at each node

Relational Model for WSNs

- Tuples include sensor readings and associated sensor types, node ID, timestamp, energy balance, etc.
- Append-only and distributed across multiple nodes, thus supporting **streamed, distributed** data
- Query is ***persistent*** and ***periodically*** evaluated
- Queries themselves are treated as data upon which other queries may operate, i.e., **recursive** query.

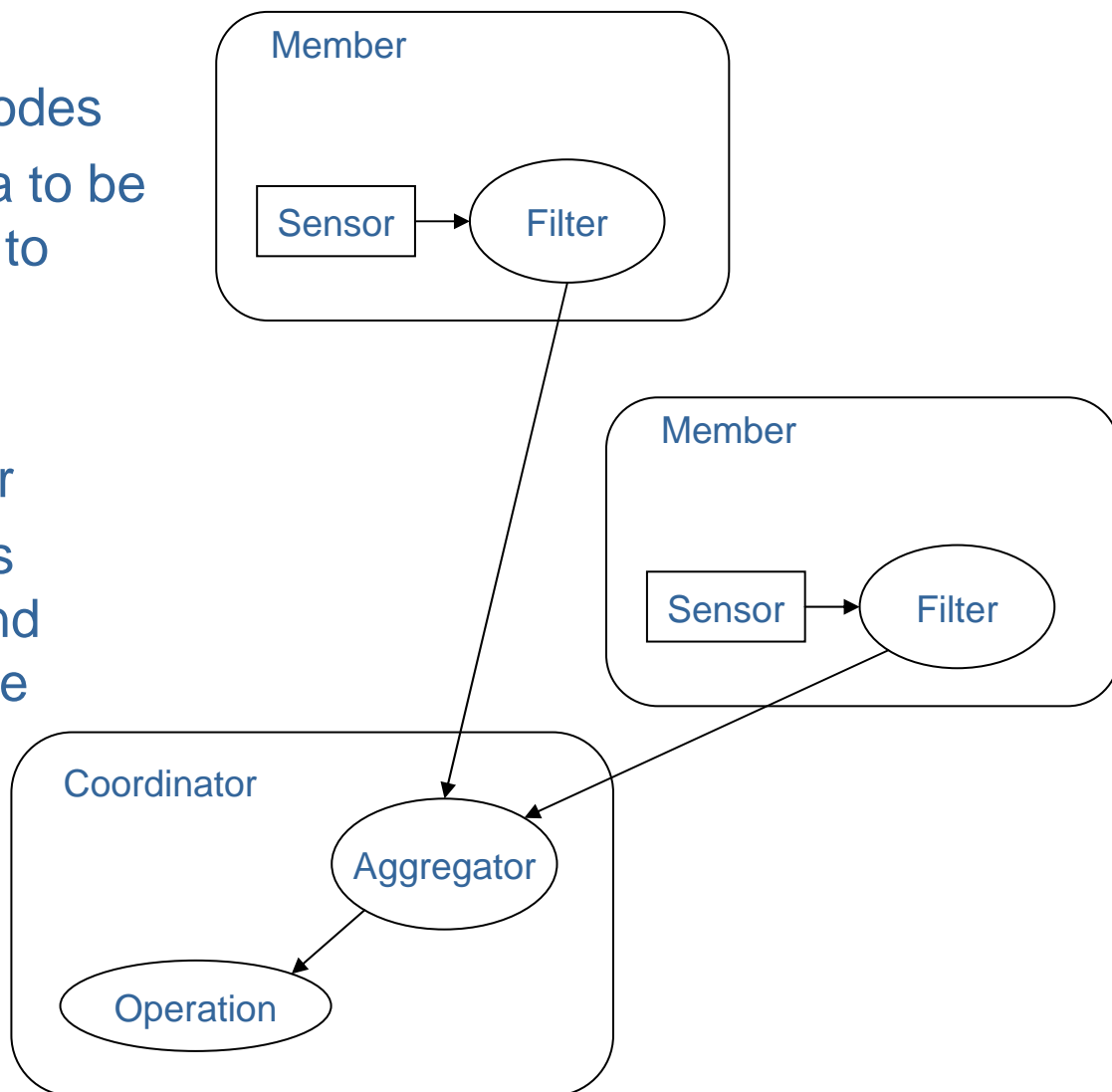
Hierarchical Architecture



- Roles
 - Super coordinator
 - Coordinator
 - Member
- Cluster
 - Nodes in a small region
 - One-hop communication
 - Redundancy
 - Sensing
 - Communication

Filters and Aggregators

- **Filters**
 - Run on member nodes
 - Determine the data to be collected and sent to aggregators
- **Aggregators**
 - Run on coordinator
 - Collect data across space and time, and perform appropriate operations



A Simple SQL-like Interface

- A sample query

```
SELECT cluster_id, AVG(mag)
```

```
FROM sensors as s
```

```
WHERE s.mag > 40
```

```
GROUP BY cluster_id
```

```
INTERVAL 1sec
```

```
DURATION 12min
```

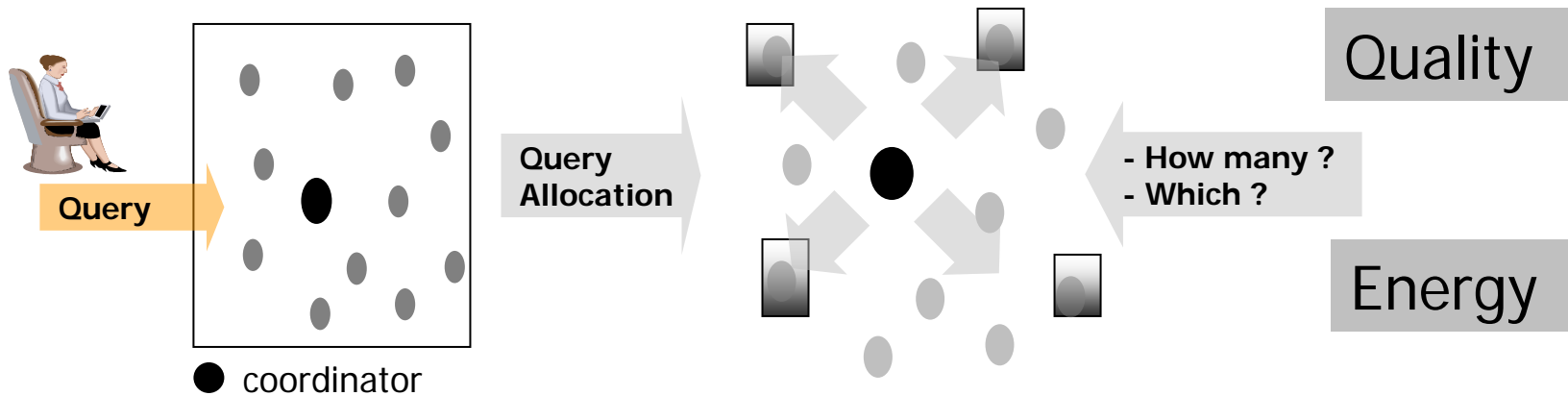
Aggregator

Filter

- Queries that operate on queries

- Insert, Delete, Update, Select, and Estimate

Energy-aware AQP

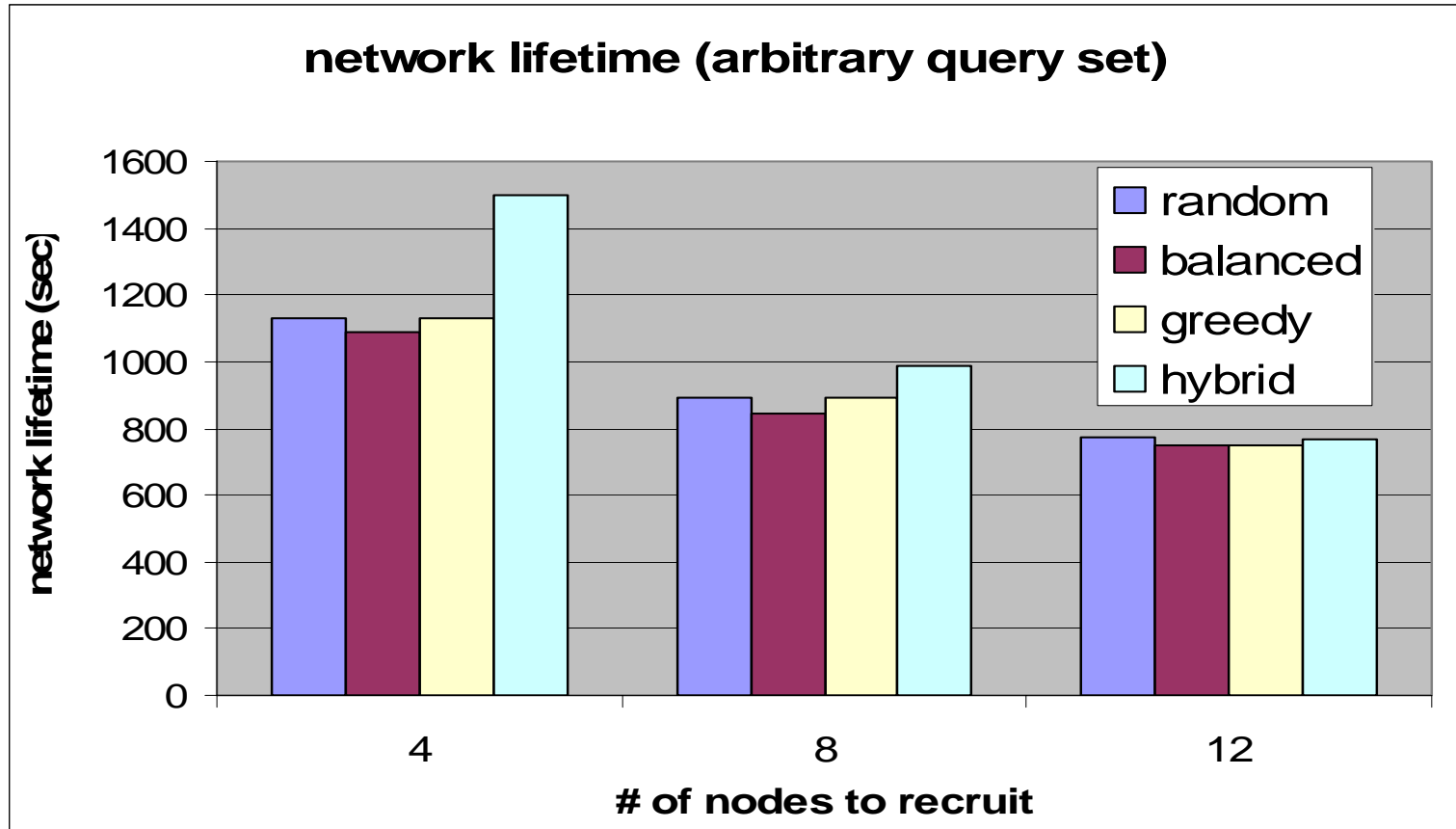


- Distribute workload using **utility/cost** model
- Given a local cluster of n substitutable nodes, adaptively distribute workload to a subset of the nodes
- **Utility**: accuracy of the query result
 - More nodes give better estimate of sensor value
- **Cost**
 - Cost associated with selecting and aggregating data
 - Models: balanced, greedy, hybrid

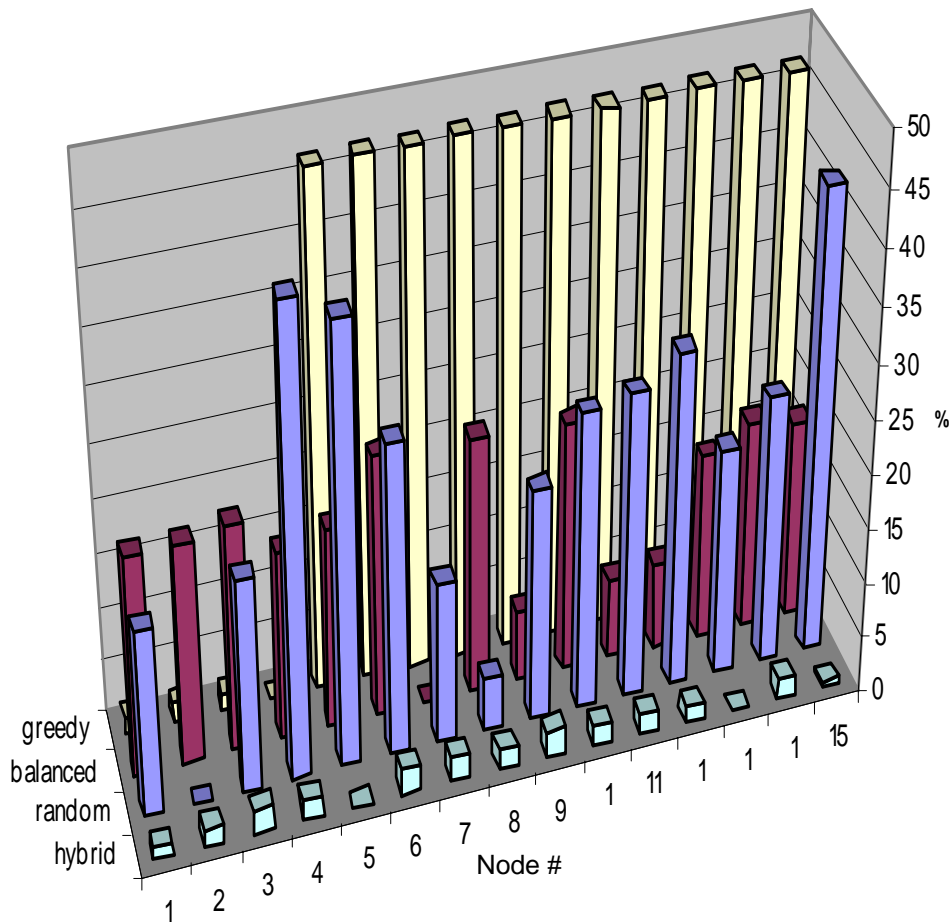
Comparison of Cost Models

Cost Model	Description
Balanced	Cost = $1/(\text{Residual Energy})$ <u>Balances</u> nodes' energy consumption
Greedy	Cost = Additional Energy Consumption <u>Minimize</u> energy consumption by adding a new query
Hybrid	A combination of Greedy+Balanced <u>Greedy</u> to allocate incoming queries and <u>Balanced</u> to exchange existing query sets

Network Lifetime



Per-node Residual Energy



- Selects four nodes per query out of 15 possible
- Remaining energy is measured at the end of network lifetime
- Hybrid model achieves a longer lifetime by distributing power usage more evenly over available nodes

Online Query Optimization

- Why?
 - Queries may be submitted at any time
 - Availability of sensor nodes may change
- Main focus of query optimization is to save energy
 - Maximize **sharing** of communication and sensing costs among queries

AQP Demonstration

- Implemented support for the “Pursuer-Evader Game” scenario
 - Tracks an enemy evader through a field
 - *Location estimate* is used to pursue the evader
- Steps
 - Energy-aware **Coordinator** election
 - Energy-aware, geographically-distributed **Sentry** assignment
 - Detection and aggregation for **estimation**
 - Adaptive estimation
 - **Re-election** of Coordinators and Sentries

Content-Aware Metadata Creation and Access

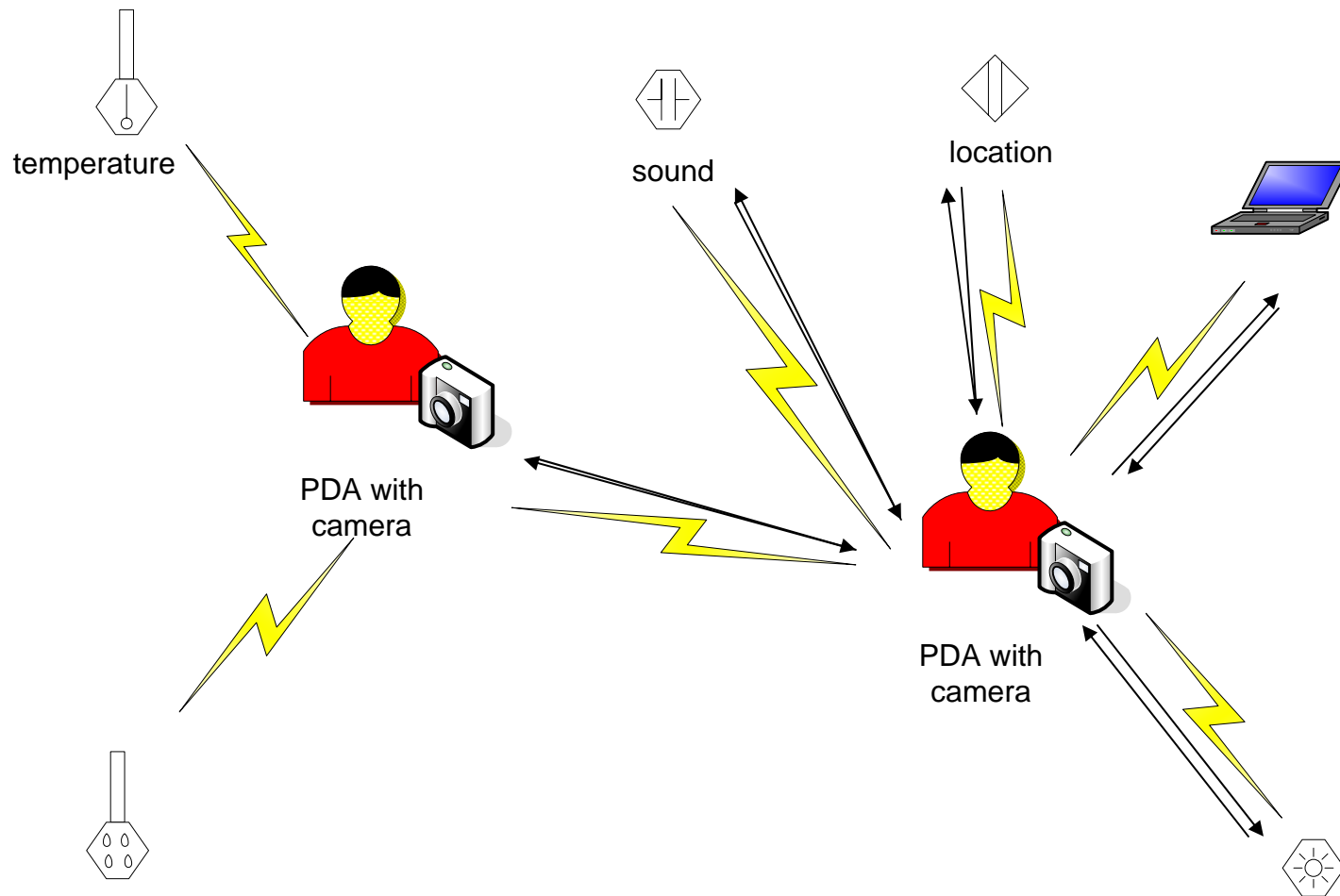
- Wireless handheld devices and sensors are becoming **everywhere!**
- Amount of digital media data is rapidly **increasing** and becoming **burdensome** to manage
 - =>Difficult to find, edit, share, and reuse media because computers don't understand media content
 - Media is opaque and data-rich and lacks structured representations

So, We

Designed a framework to:

- **Collect** environmental information from wirelessly-enabled devices
- **Associate** the collected information, or “metadata,” with digital media files
- Metadata **facilitates** easy search, categorization, and organization of files.

Communication Model



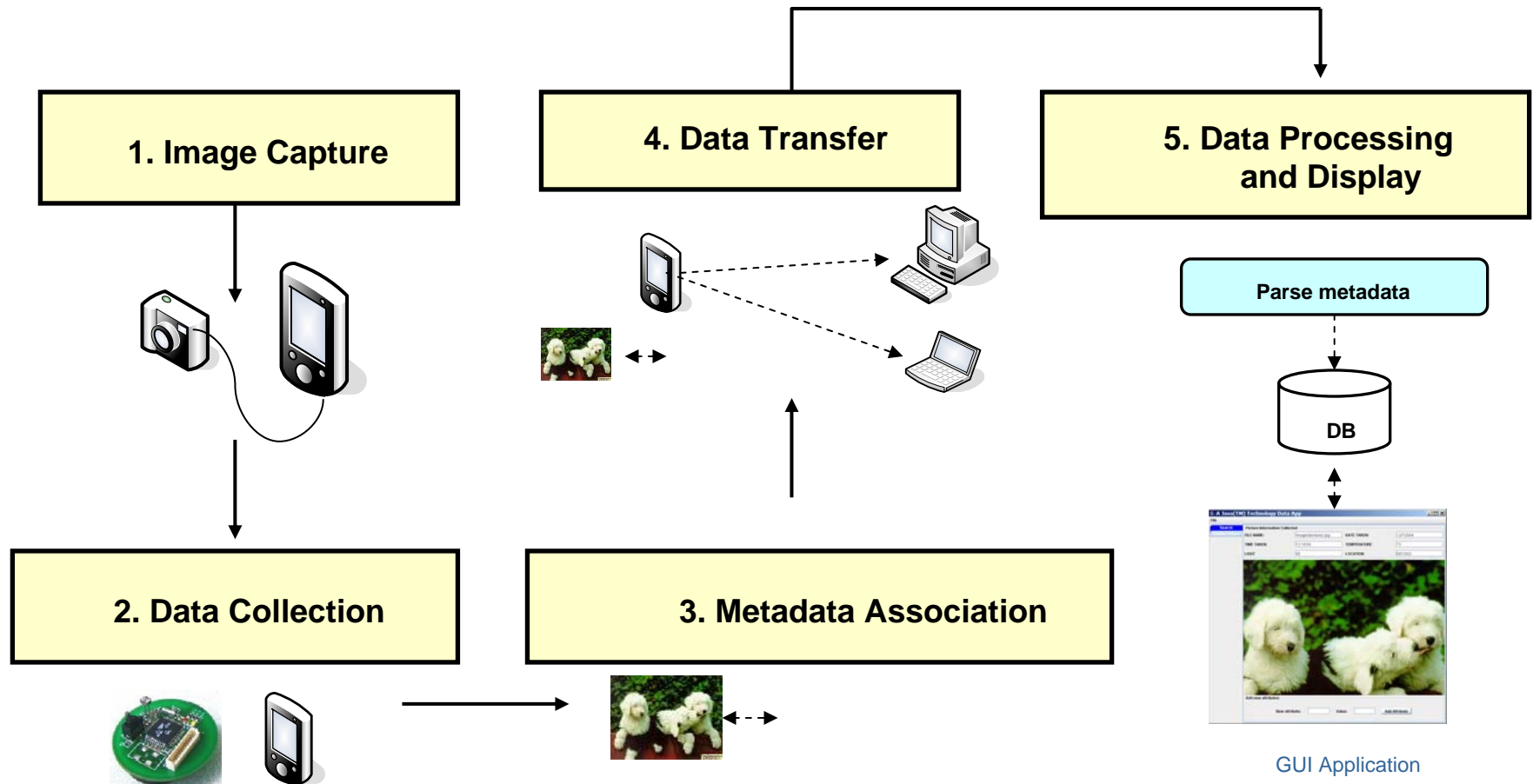
Heterogeneous Networks

- **Mobile users** (iPAQs & Stargates)
 - User input simulates taking pictures
 - 802.11 WLAN communication
- **Environmental sensors** (motes & RFIDs)
 - Measure temperature, light, and location
 - RF communication
- **Logical sensors** (laptops quipped with motes/RFIDs)
 - Communicate with mobile users and environmental sensors
 - 802.11 WLAN communication
 - RF & Bluetooth communication

Metadata Association

```
procedure Metadata_Association
{
  mark photo-shoot time;
  wait 1 association period after photo;
  determine relevant time interval;
  associate file name and timestamp;
  while ( Pop the smallest offset Data
          within relevant time interval )
    if ( !duplicated (Data) && !filtered (Data) )
      write Data to metadata;
}
```

Context-Aware Image Creation



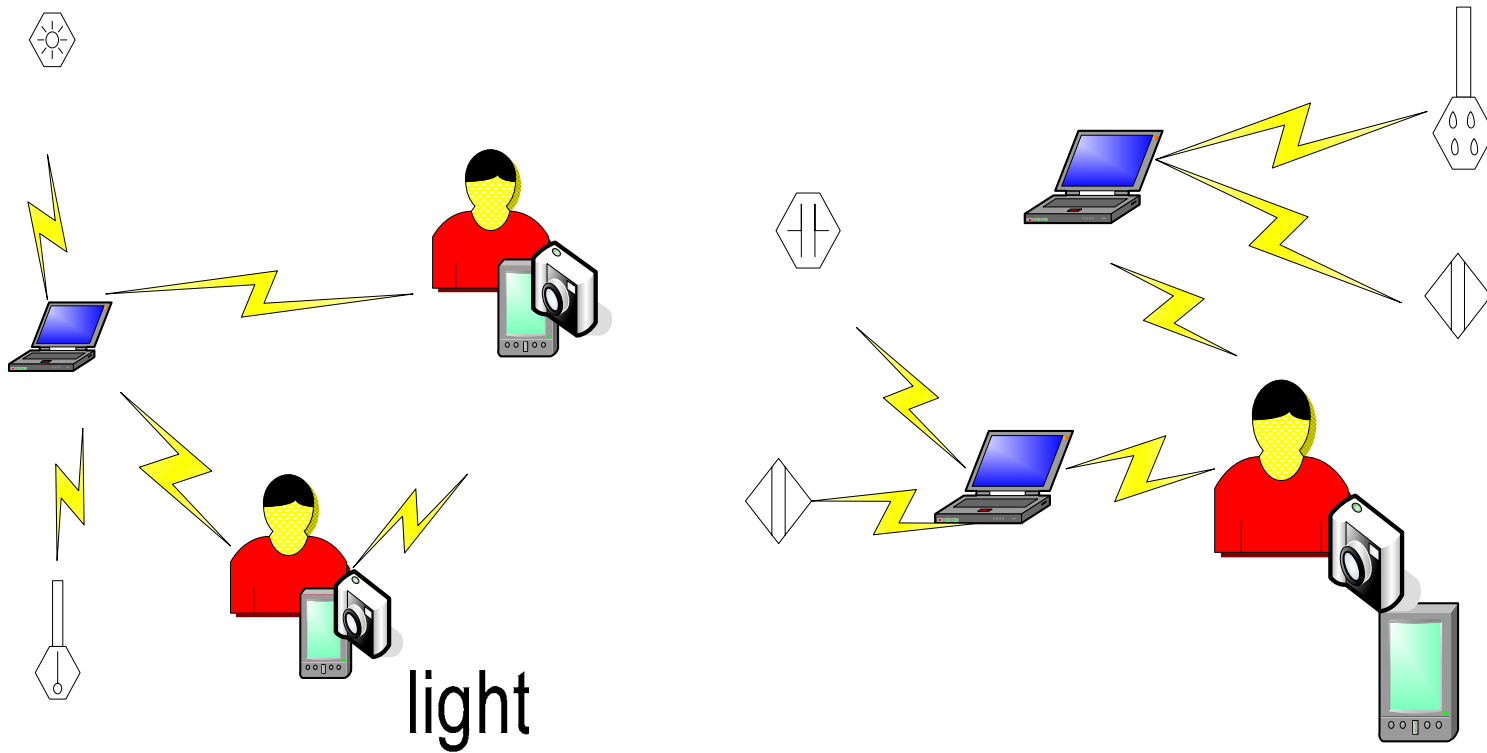
Database and GUI

- Images and associated metadata are transferred to a desk/lap-top PC server
- XML parsed and loaded into the database
- GUI application allows for flexible search and edit

Prototyping and Experimentation

- **Testbed**
 - 2 mobile nodes (iPAQs)
 - 3 logical nodes
 - 13 environmental sensors
- Users **walk** around, **take** pictures, and **collect** environmental data
 - 1-hour simulation
 - Two users at a time, total of 9 users
- **Data collection**
 - On-demand
 - Periodic

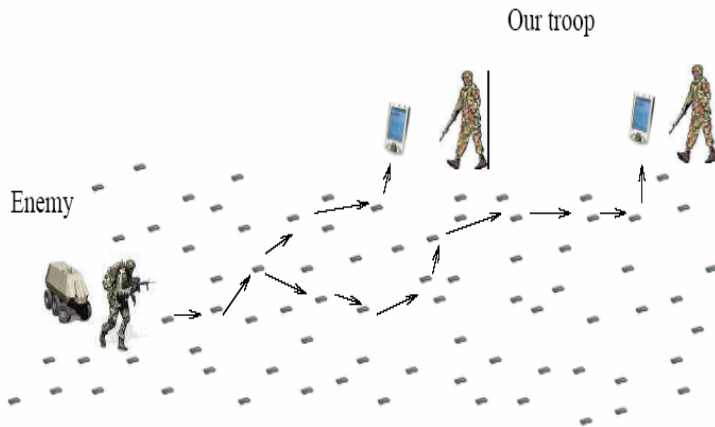
Experimental Setup



Distributed Location Service

A Typical Scenario:

- Mobile nodes issue queries to the “static” sensor network
- Query results are returned to the requester mobiles



When query results are generated:

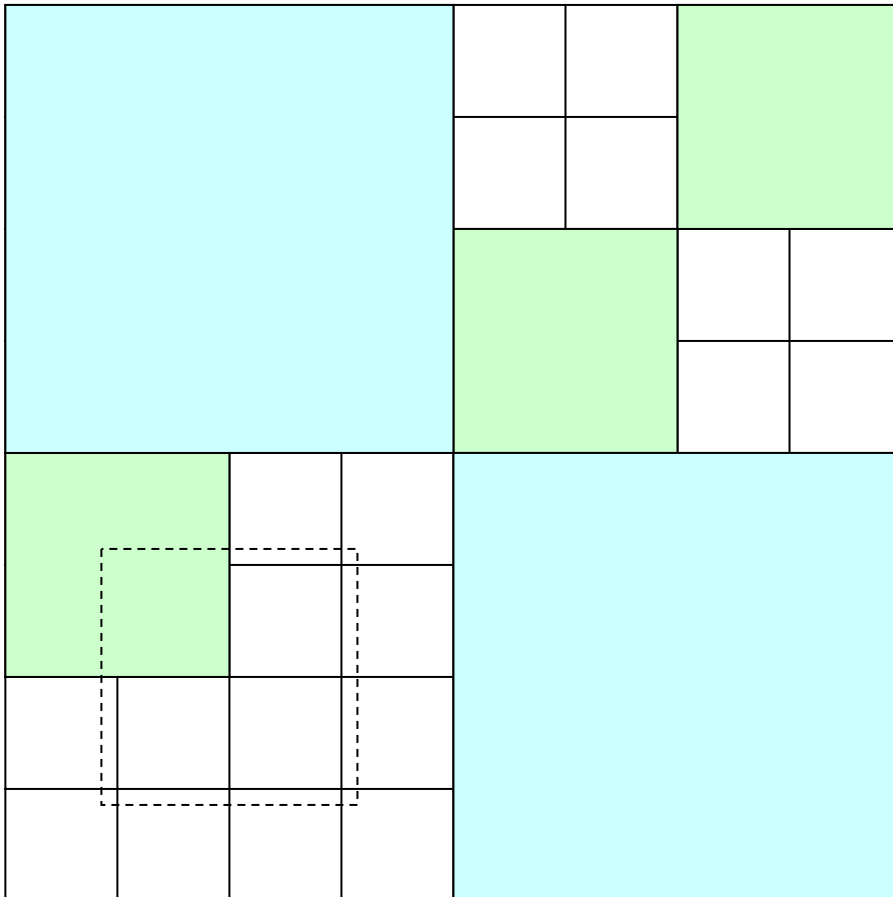
- Mobile nodes which issued query may have moved away
- Need to route sensed data to a **mobile** sink!

DLSP: Distributed Location Service Protocol

- **What does it do?**
 - provides the updated location information of mobile sinks to static sensor nodes
- **How?**
 - Each mobile independently elects location servers
 - Location info of mobiles is sent to their location servers
 - Other nodes contact the location servers to obtain the location of mobile sinks

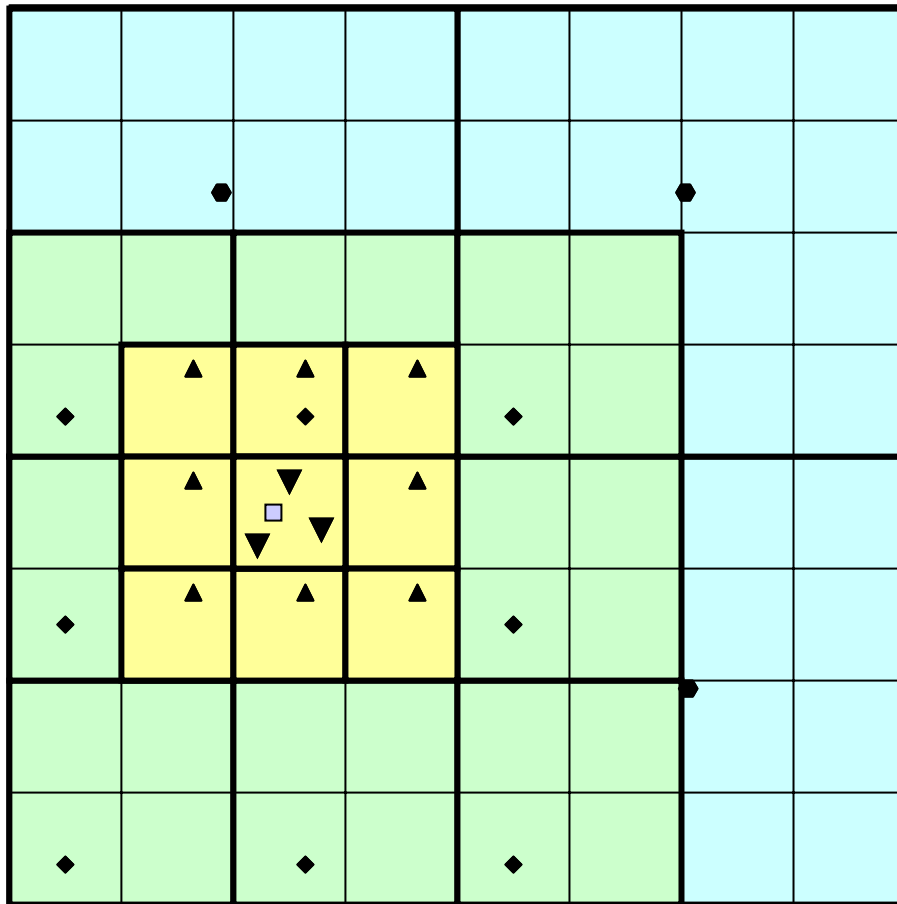
Grid Construction

- **Level-1 Square**
 - Smallest square
- **Level-(k+1) Square**
 - 4 of level-k square
- No overlap between squares of the same level
- Squares at each level cover the entire coverage area



- Level-1 square ■ Level-2 square
- Level-3 square □ Not a Level-2 square

Location Server Election



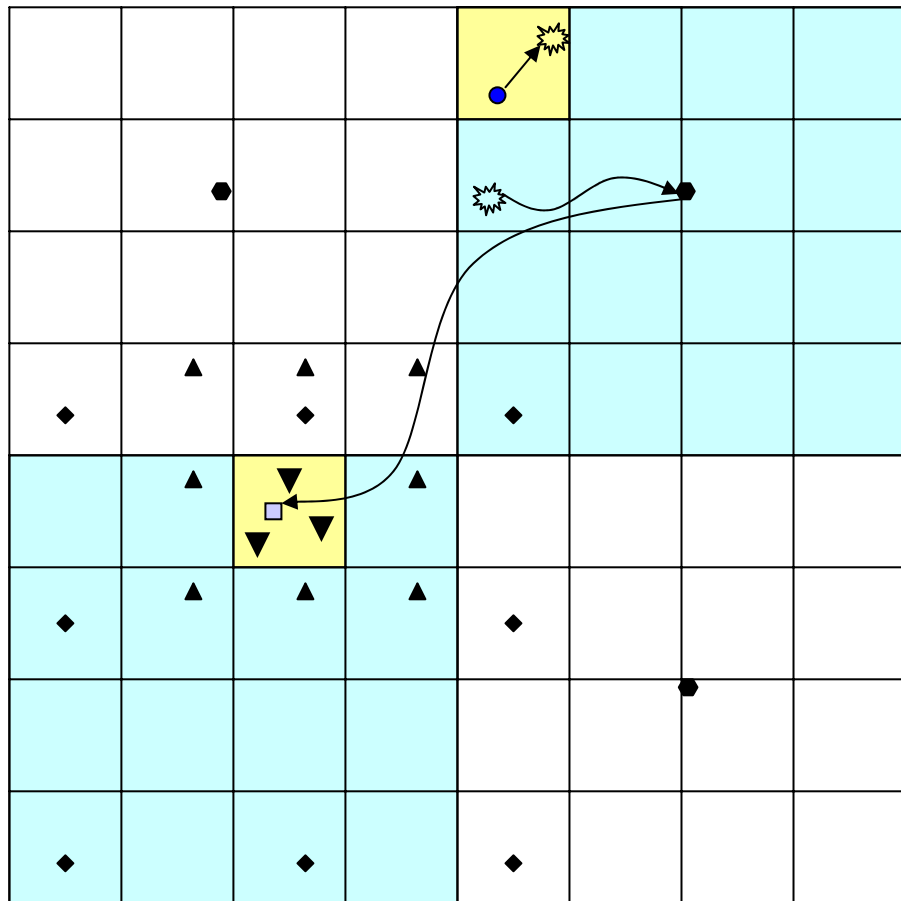
- **Level-0 Servers**
 - All the nodes within the same level-1 square
- **Level-k Servers**
 - One from each of neighboring level-k squares
 - Relative location: $H(id,k)$
- **Denser** near M and **sparser** away from M

□ Mobile Node M

▼ $DLS_0(M)$ ▲ $DLS_1(M)$

◆ $DLS_2(M)$ ● $DLS_3(M)$

Location Query



- Sink node issues a query if it needs the location of M
- Query is recursively passed to the higher-level (presumed) server

- Mobile Node M ● Source node
- ▼ $DLS_0(M)$ ▲ $DLS_1(M)$
- ◆ $DLS_2(M)$ ● $DLS_3(M)$

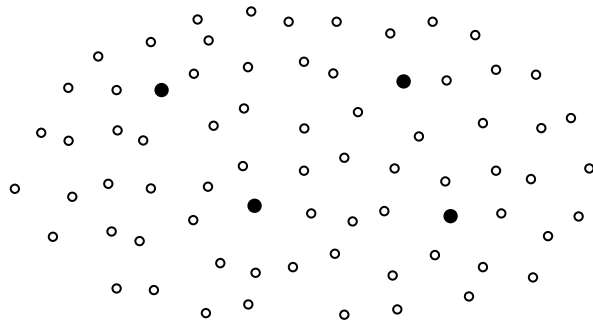
Overhead of DLSP

- **Location Query**
 - d : distance between src and dst
 - # of msg/query: $O(d)$
 - delay/query: $O(d)$
- **Location Information Maintenance**
 - N : # of sensor nodes, M : # of mobile nodes, L : network size (distance)
 - Mem requirement per sensor node: $O(M \cdot \log(N)/N)$
 - # of msg/mobile node/period: $O(L \cdot \log(N))$

Comparison with Others

- **MIT's GLS**
 - GLS: Every node is assumed mobile
 - DLSP: Only a small portion of nodes are mobile
=> more efficient
- **Landmark routing**
 - DLSP: No need to maintain landmark hierarchy
(when nodes move, die, etc.)
- **TTDD**
 - No overhead for query forwarding, double agent, and local query re-flooding

Security in Networked Embedded Systems



Sensor Network

- Self-organizing, self-healing
- Battery-powered
- Unattended, not rechargeable
- A large number of nodes

CHALLENGES

- Wireless
- Limited Energy
- Large-scale



OUR APPROACH

- Lightweight
Not sacrificing security level
- Distributed, P2P
- Tailored to Threat/Svc

Threat Model

OUTSIDER

Data Attacks

- Traffic capture/replay
- Spoofing if unencrypted
- Man-in-the-middle (limited)

Radio Attacks

- High-power jamming
- Radio source detection

Physical Attacks

- Reprogram as malicious
- Destroy device
- Extract key materials

INSIDER

Data Attacks

- Traffic injection/flooding
- Unlimited spoofing
- DoS, Man-in-the-middle

Service Disruption on

- Routing (altered/selective)
- Clock synchronization
- Localization

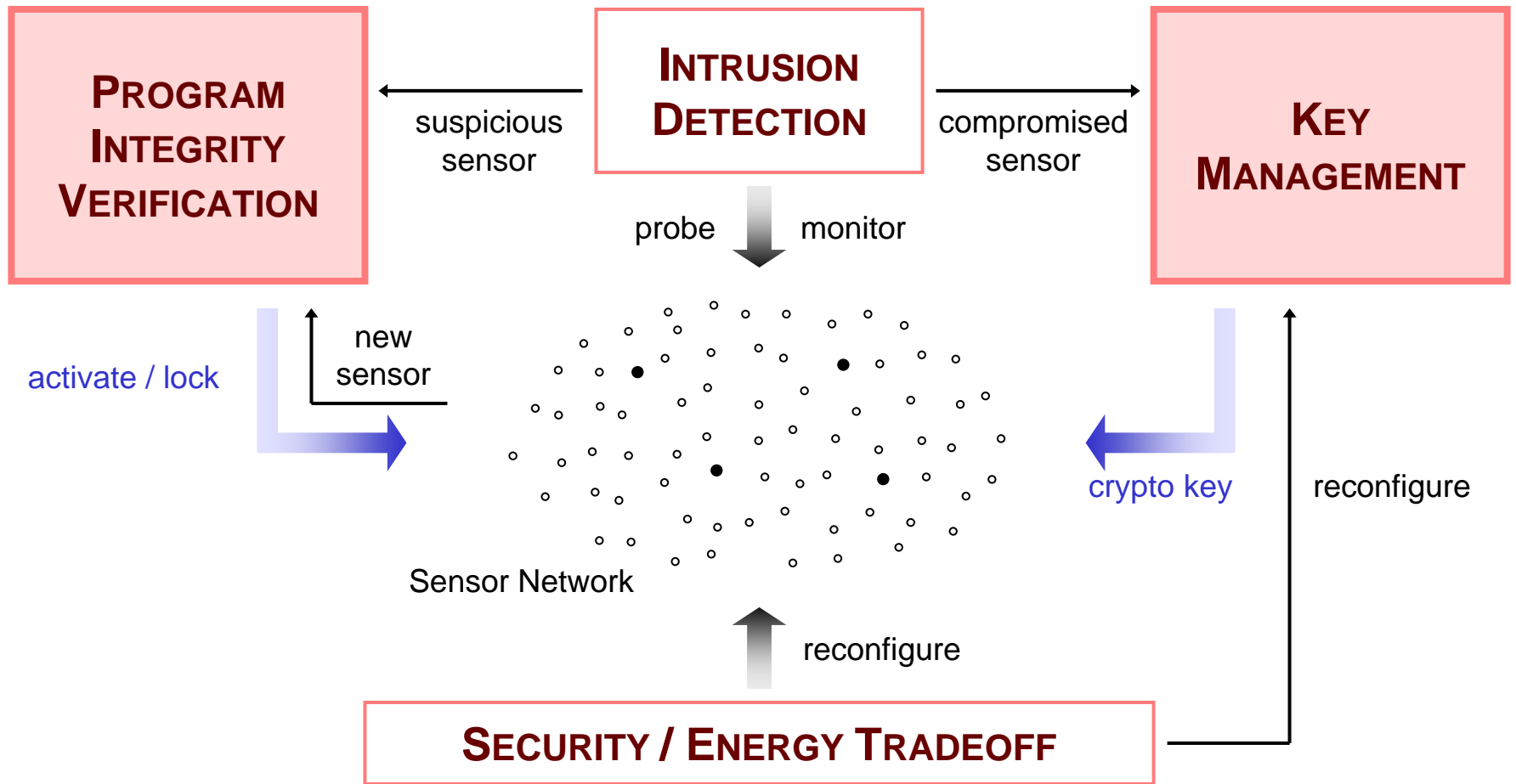
Miscellaneous

- Service/data to adversary
- Malicious service to net

Why LiSP?

THREAT	DEFENSE	PROBLEM	SOLUTION
Attack on Traffic <ul style="list-style-type: none">EavesdroppingTraffic replay, modification, injectionService disruption, DoS	Key Sharing <ul style="list-style-type: none">GloballyGroup-basedPairwise Re-Keying <ul style="list-style-type: none">PeriodicallyEvent-triggered	<ul style="list-style-type: none">Vulnerable to sensor compromisesLarge re-keying overheadTranscoding per hop	Group-based Key Management Two-Tier Nets Distributed Key Management P2P Nets
Attack on Program <p>The adversary can</p> <ul style="list-style-type: none">capturereverse-engineerre-programclone sensor device(s)	H/W Tamper-Resistance S/W <ul style="list-style-type: none">ObfuscationResult CheckingSelf-Decryption	Protection of program itself → Defenseless once broken	Soft Tamper-Proofing via Program-Integrity Verification

LiSP Architecture

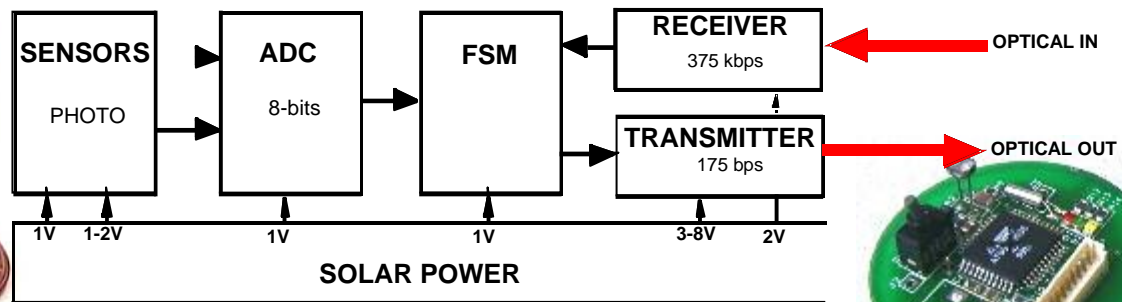
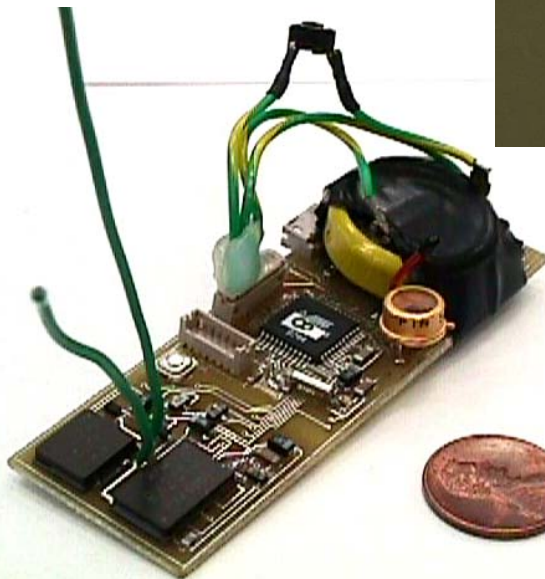
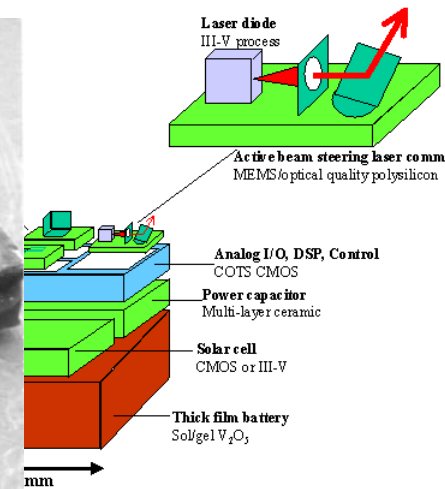
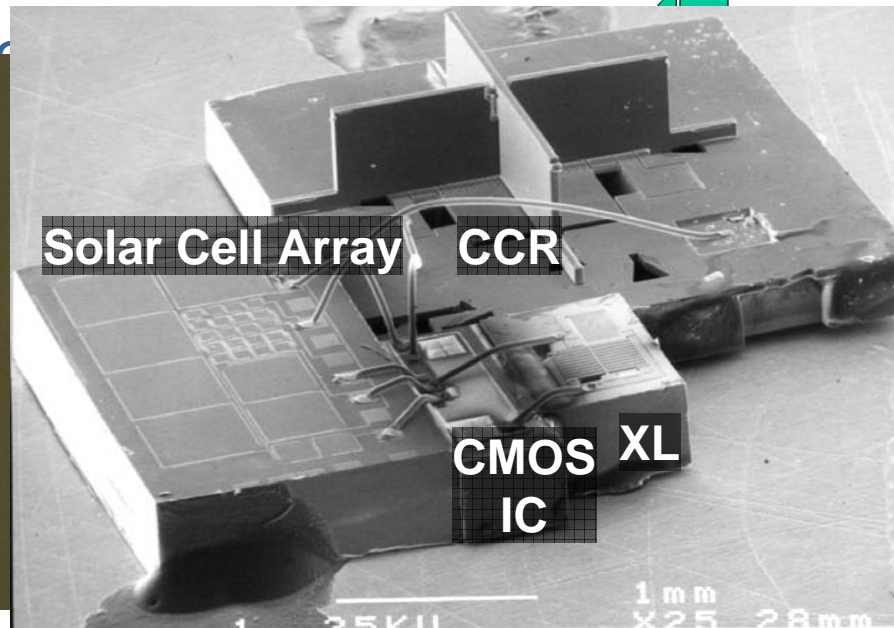


Sensor Networks Research at UCB

Miniaturization – Pister (SmartDust)

- The Goal: Autonomous millimeter-scale robot
 - Sensing
 - Computation
 - Communication
 - Power
 - Motors

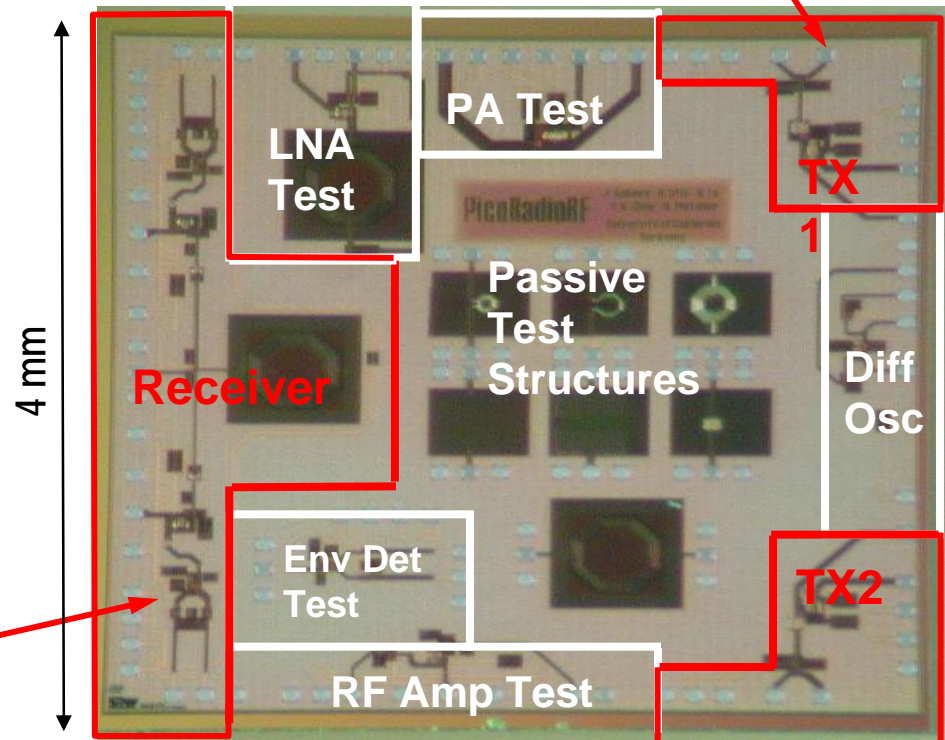
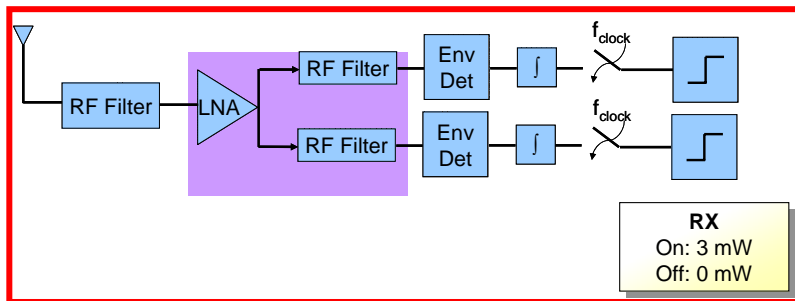
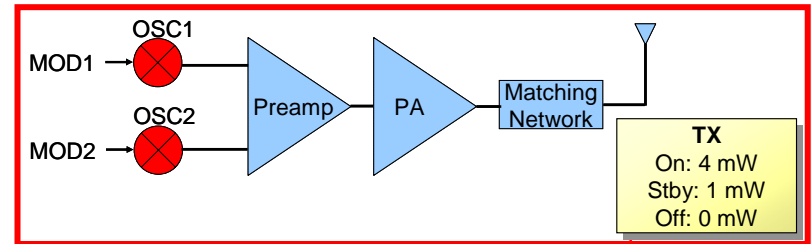
Smart Dust Components



Low Power RF – Rabaey (PicoRadio)

- CMOS
 - Cheap, Integrated
- mW -> sub mW
- Simple
- Advantage in Numbers

BWRC



System/Networking/Programming – Culler

Services

Networking

TinyOS

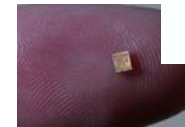
www.tinyos.net

WeC 99
"Smart Rock"

Rene 11/00

Dot 9/01

Mica 1/02



Small microcontroller

- 8 kb code,
- 512 B data

Simple, low-power radio

- 10 kb

EEPROM (32 KB)

Simple sensors

Designed for experimentation

- sensor boards
- power boards

DARPA SENSIT, - Intel Expeditions

Crossbow

Demonstrate scale

NEST open exp. platform

128 KB code, 4 KB data

50 KB radio

512 KB Flash

comm accelerators

- DARPA NEST

Structural Monitoring – Glaser, Fenves

- Dense Instrumentation of Full Structure
 - Cost is all in the wires
- Leads to in situ monitoring

Liquifaction, Tokashi Port



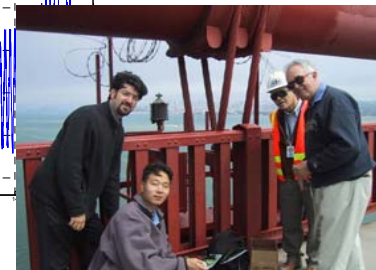
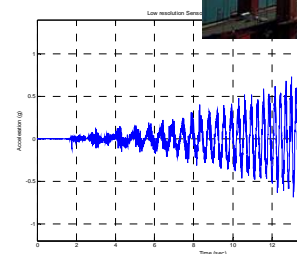
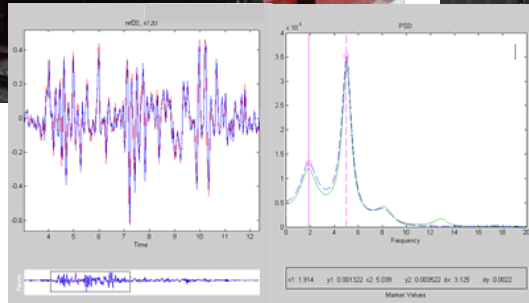
25 Motes on Damaged sidewall



30 Motes on Glue-lam beam



Wind Response Of Golden Gate Bridge



Protection – Sastry, Culler, Brewer, Wagner



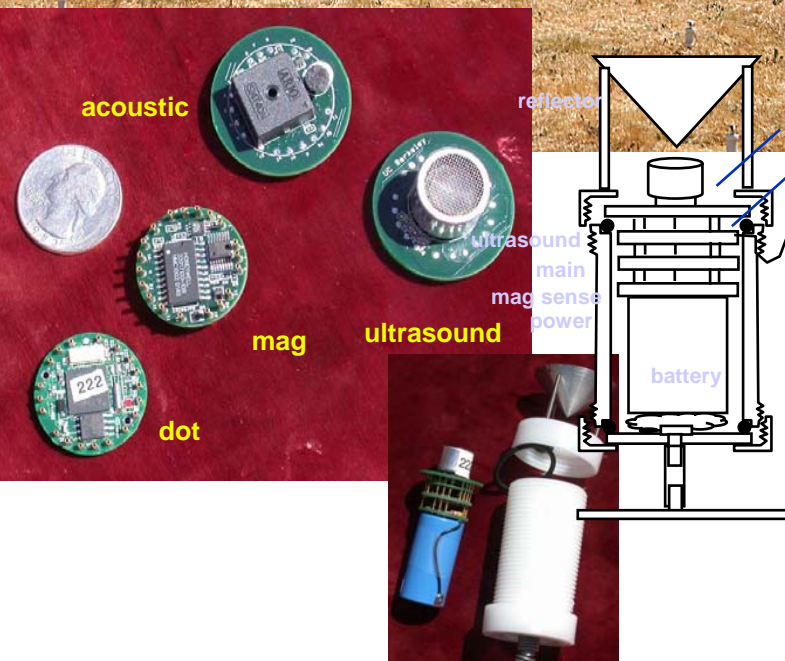
Detect vehicle entering sensitive area, track using magnetics, pursue and capture by UGV.

Components

- 10x10 array of robust wireless, self-localizing sensors over 400 m² area
- Low cost, robust 'mote' device
- Evader: human controlled Rover
- Pursuer: autonomous rover with mote, embedded PC, GPS

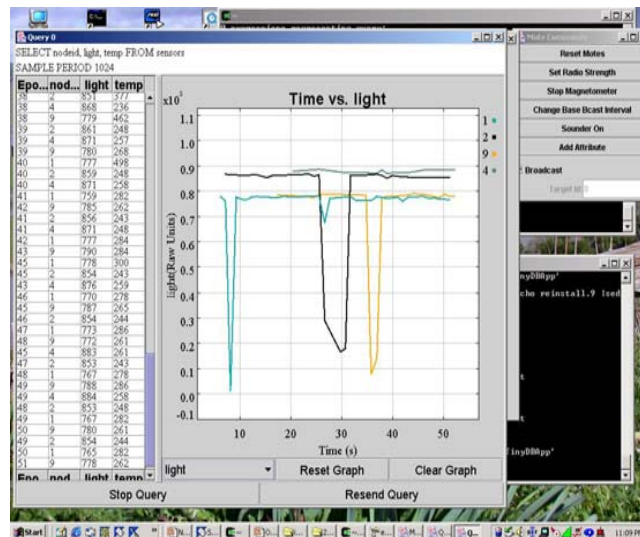
- Operation

- Nodes inter-range (Ultrasonic) and self localize from few anchors, correct for earth mag, go into low-power 'sentry' state
- Detect entry and track evader
 - Local mag signal processing determines event and announces to neighbors
 - Neighborhood aggregates and estimates position
 - Network routes estimate from leader to tracker (multihop)
- Pursuer enters and navigates to intercede
 - Motes detect and estimate multiple events
 - Route to mobile Pursuer node
 - Disambiguates events to form map
 - Closed inner-loop navigation control
 - Closed information-driven pursuit control

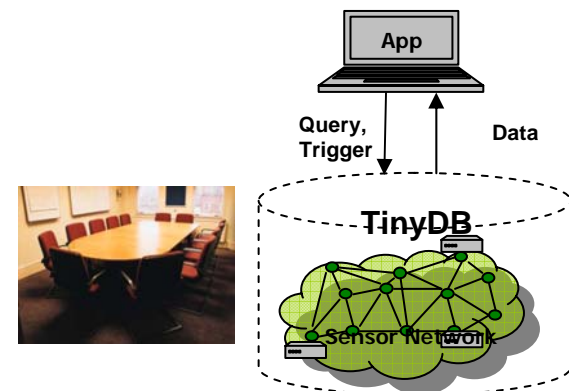


Sensor Net Databases – Hellerstein, Franklin

- Relational databases: rich queries described by declarative queries over tables of data
 - select, join, count, sum, ...
 - user dictates **what** should be computed
 - query optimizer determines **how**
 - assumes data presented in complete, tabular form
- database operations over streams of data
 - incremental query processing
- process the query **in** the sensor net
 - query processing == content-based routing?
 - energy savings, bandwidth, reliability

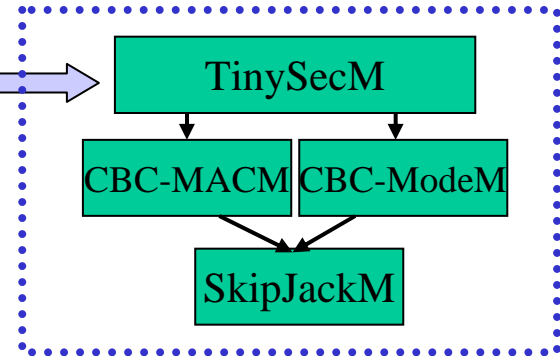


`SELECT AVG(light)`
`GROUP BY roomNo`



Security - Wagner

cellphones



**Let's get it right
the first time!**

wireless networks

sensor networks

1980 analog cellphones: AMPS

analog cloning, scanners
▶ fraud pervasive & costly

digital: TDMA, GSM

TDMA eavesdropping [Bar]

more TDMA flaws [WSK]
GSM cloneable [BGW],
GSM eavesdropping [BSW,BGW]

2000 Future: 3rd gen.: 3GPP, ...

1999 802.11, WEP

2000

2001 WEP broken [BGW]
WEP badly broken [FMS]
▶ attacks pervasive

2002

2003 WPA

Future: 802.11i

Berkeley notes

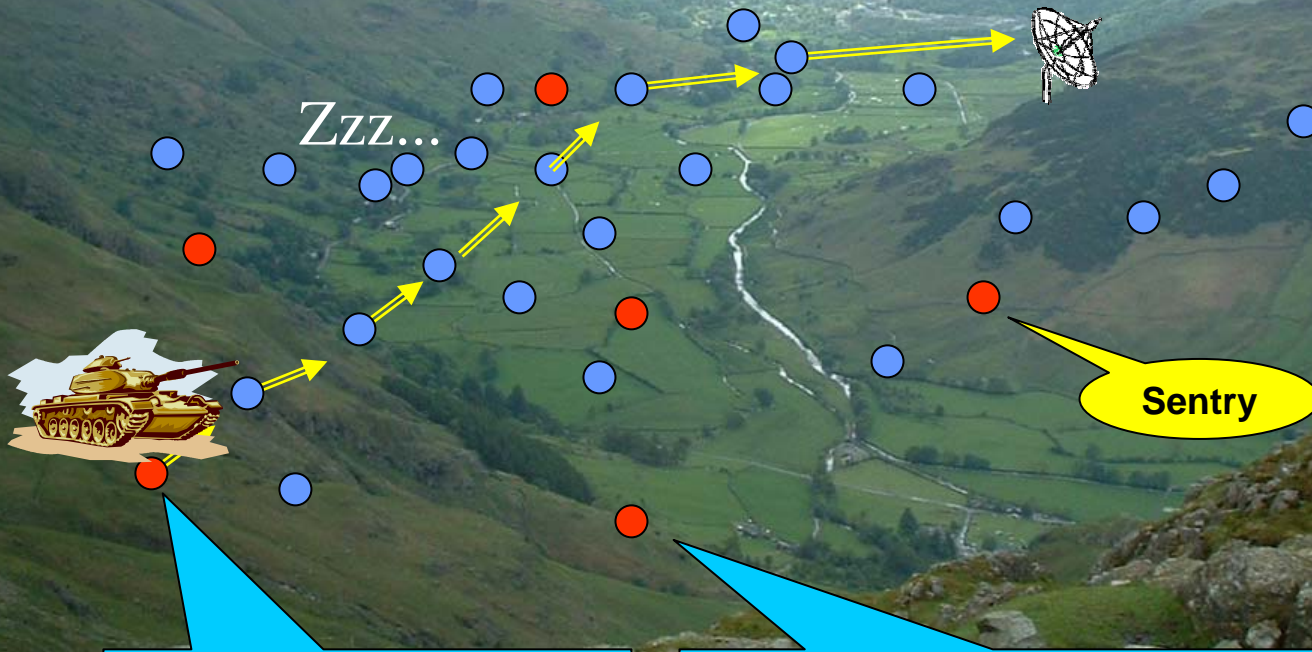
2002 TinyOS 1.0

2003 TinyOS 1.1, TinySec

VigilNet
University of Virginia

Energy Efficient Surveillance System

1. An unmanned plane (UAV) deploys motes



3. Sensor network detects vehicles and wakes up the sensor nodes

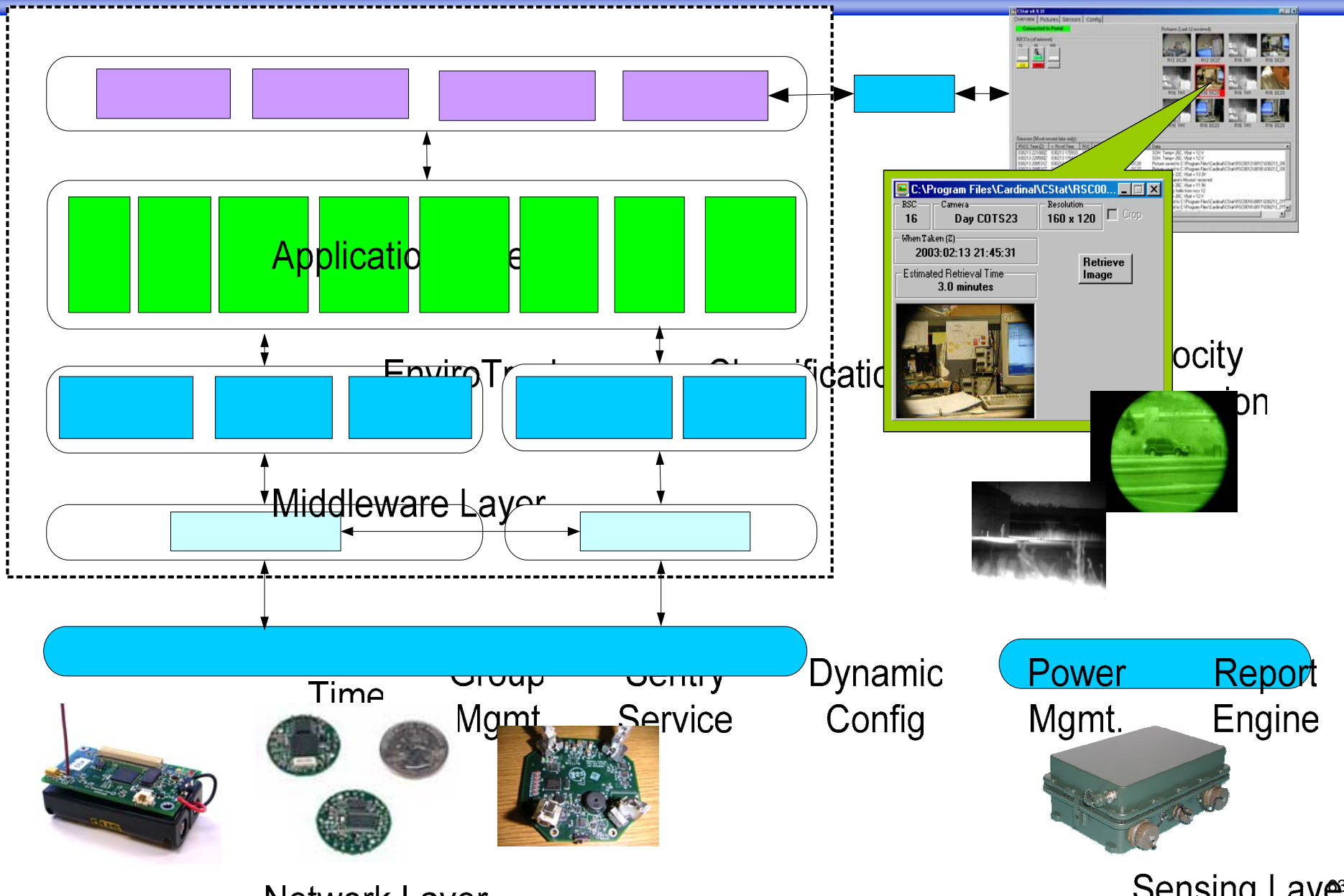
2. Motes establish an sensor network with power management

Diffusion Routing
Neighbor Discovery
Time Synchronization
Parameterization
Sentry Selection
Coordinate Grid
Data Aggregation
Data Streaming
Group Management
Leader Election
Localization
Network Monitor
Tripwire Service
Reconfiguration
Reliable MAC
Leader Migration
Scheduling
State Synchronization
.....

Goals

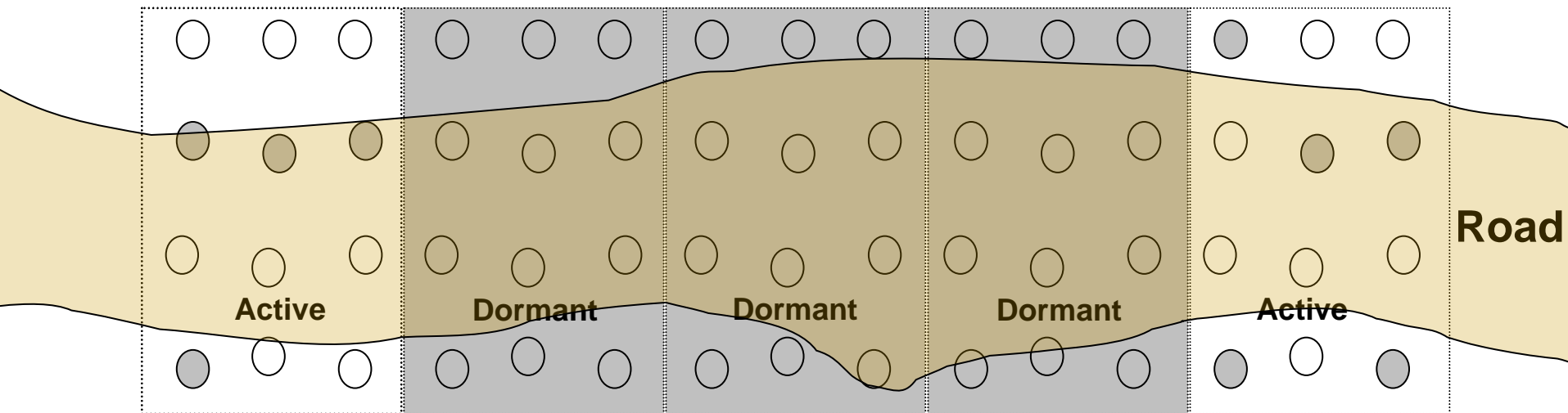
- Develop an operational self-organizing sensor network of size 1000
- Cover an area of 1000m x 100m
- Stealthy
- Lifetime 3-6 months
- Timely detection, track and classification
 - Large or small vehicle
 - Person, person with weapon
- Wakeup other devices when necessary
 - Extend the lifetime of those devices as well
- Exhibit self-healing capabilities

VigilNet Architecture V1.3

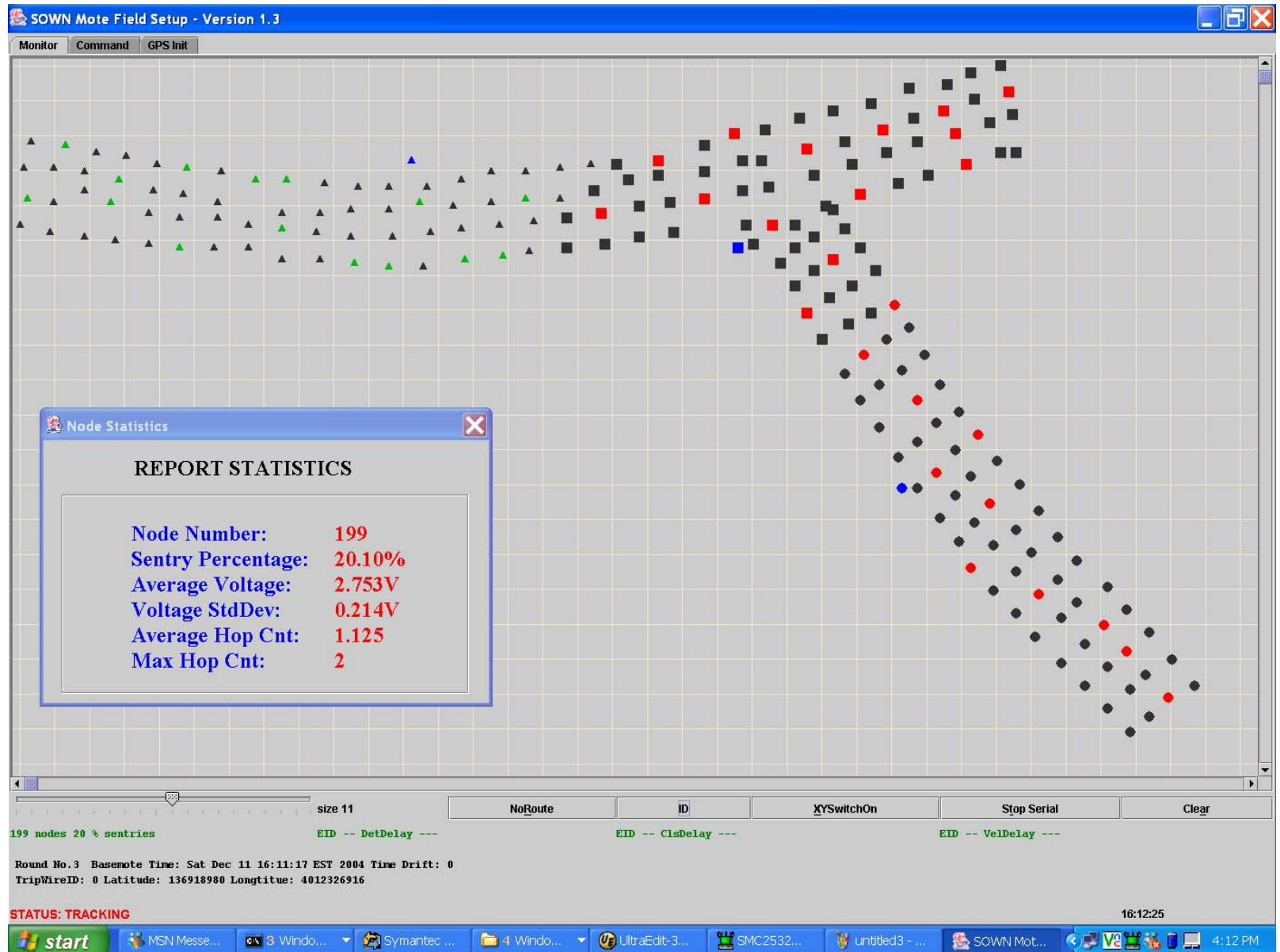


Tripwire-based Surveillance

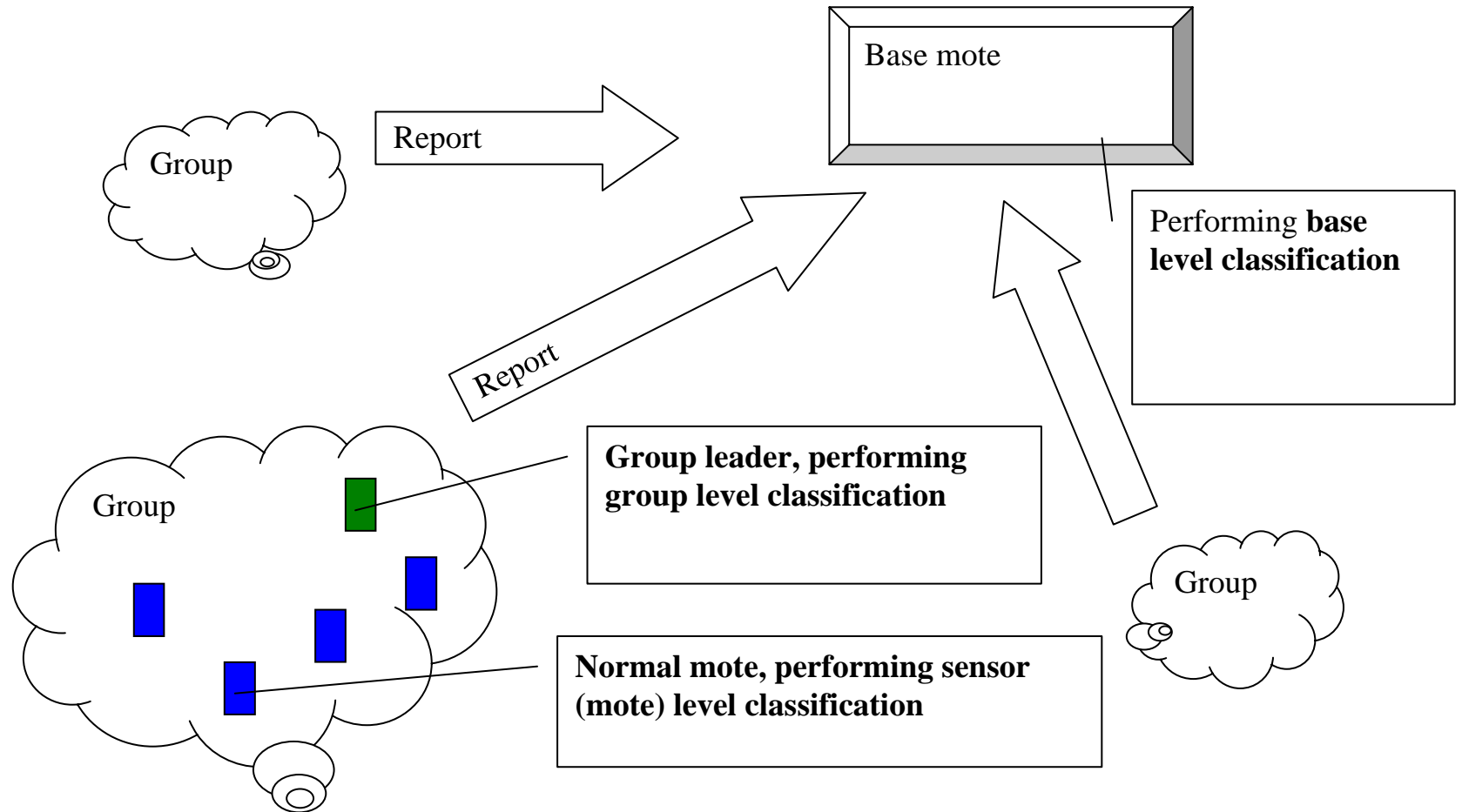
- Partition sensor network into multiple sections.
- Turn off all the nodes in dormant sections.
- Apply sentry-based power management in tripwire sections
- Periodically, sections rotate to balance energy.



System Test with 203 Nodes



3-Tier Classification



Concluding Remarks

- Sensor networks provide an inexpensive vehicle for exploring various (old and new) research issues
- Commercial applications with RFIDs as leader
- Current and future directions: query processing using geostatistics, sensor network security; tradeoffs among perf, security, reliability and resource consumption; extreme scaling and other DoD/commercial apps.